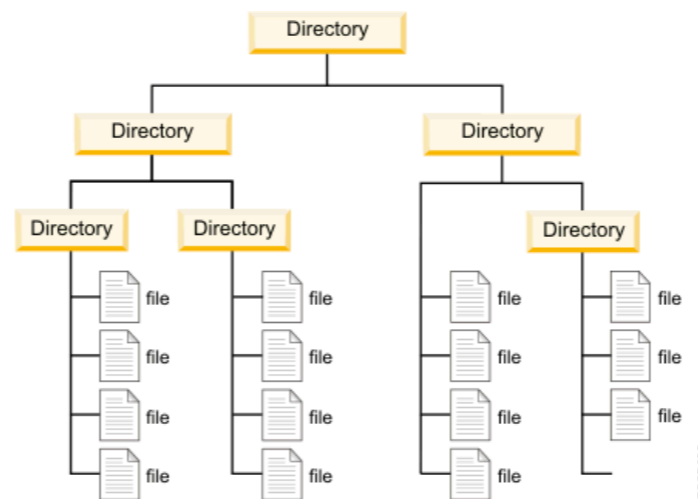
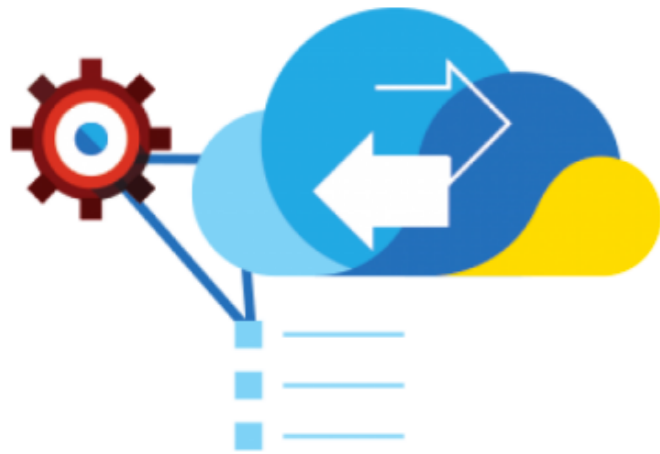


Synthesizing Bijective String Transformers

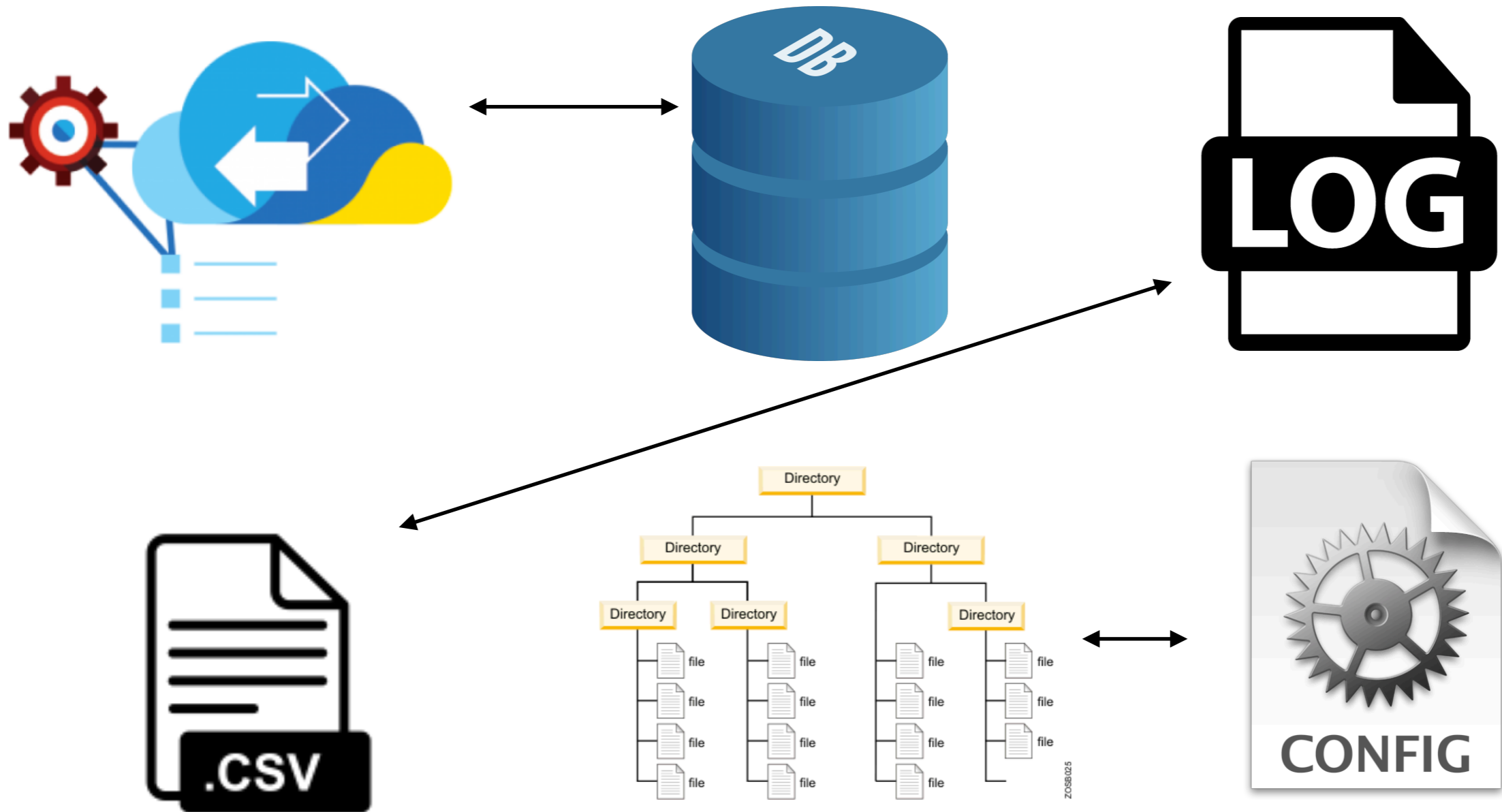
Anders Miltner
David Walker (advisor)
Kathleen Fisher
Benjamin Pierce
Steve Zdancewic



Data is Everywhere



Data is Everywhere



Example: BibTex and Endnote

BibTex

```
@book{
  title={Introduction to
    MetaMathematics},
  author={Kleane, Stephen
    Cole},
}
```

EndNote

```
%0 Book
%T Introduction to
  MetaMathematics
%A Stephen Cole Kleane
```

Example: BibTex and Endnote

BibTex

```
@book{
  title={Introduction to
        MetaMathematics},
  author={Kleene, Stephen
        Cole},
}
```

EndNote

```
%0 Book
%T Introduction to
  MetaMathematics
%A Stephen Cole Kleane
```

Example: BibTex and Endnote

BibTex

```
@book{
  title={Introduction to
        MetaMathematics},
  author={Kleene, Stephen
        Cole},
}
```

EndNote

```
%0 Book
%T Introduction to
  MetaMathematics
%A Stephen Cole Kleene
```

Example: BibTex and Endnote

BibTex

```
@book{
  title={Introduction to
        MetaMathematics},
  author={Kleene, Stephen
        Cole},
}
```

EndNote

```
%0 Book
%T Introduction to
  MetaMathematics
%A Stephen Cole Kleene
```

Example: BibTex and Endnote

BibTex

```
@book{
  title={Introduction to
    MetaMathematics},
  author={Kleene, Stephen
    Cole},
}
```

EndNote

```
%0 Book
%T Introduction to
  MetaMathematics
%A Stephen Cole Kleene

%0 Article
%T Type-and-Example-Directed
  Program Synthesis
%A Peter-Michael Osera
%A Steve Zdancewic
```


Example: BibTex and Endnote

BibTex

```
@book{
  title={Introduction to
        MetaMathematics},
  author={Kleene, Stephen
        Cole},
}

@article{
  title={Type-and-Example-
        Directed Program
        Synthesis},
  author={Osera, Peter-Michael
        and Zdancewic, Steve},
}
```

EndNote

```
%0 Book
%T Introduction to
  MetaMathematics
%A Stephen Cole Kleene

%0 Article
%T Type-and-Example-Directed
  Program Synthesis
%A Peter-Michael Osera
%A Steve Zdancewic
```

Bidirectional Programming

- Use a single piece of code to express functions between data formats
- Desire code to be correct by construction
 - Inversion properties guaranteed
- We focus on a bijective fragment of Boomerang [BFPPS '08]
 - Terms are “lenses”
 - Mappings between string data formats
 - Generated functions are inverses of each other

Issues with Bidirectional Programming

- Difficult to program in!
 - Complicated Type Systems
 - New Syntax
 - Fiddly details
 - New way of thinking about programs

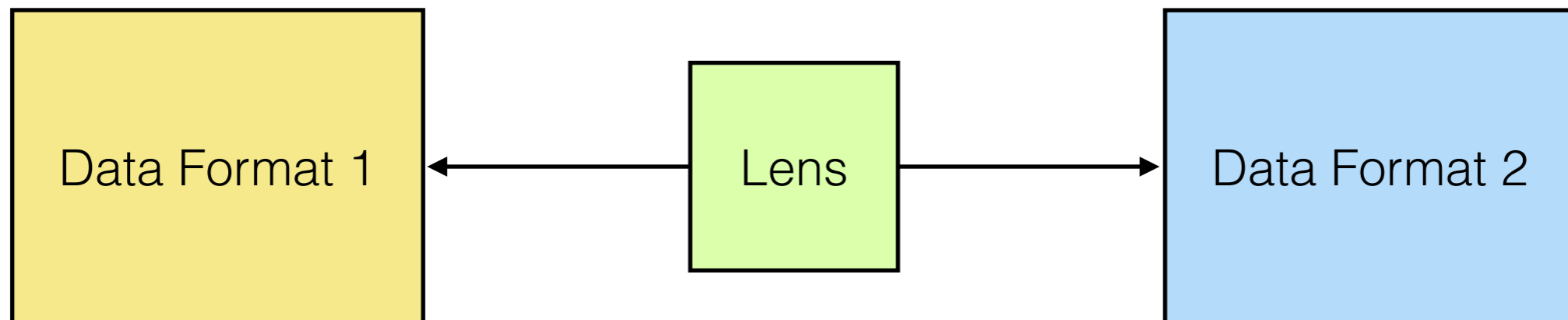
Issues with Bidirectional Programming

Miltner, Anders, NJ, \
08544, Princeton, 11 Dickinson \
St., Apt 11

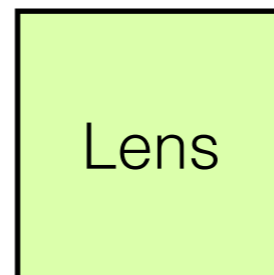
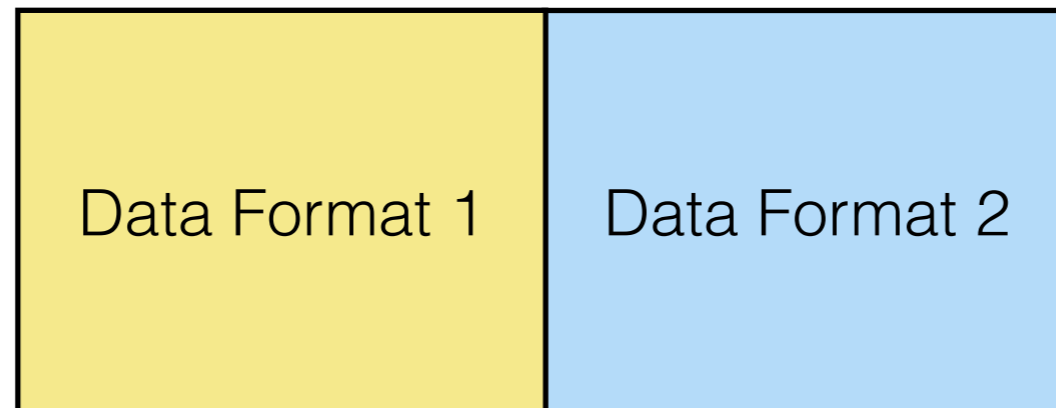
Anders Miltner
11 Dickinson St.
Apt 11
Princeton, NJ 08544

```
concat(swap(concat(...),swap(swap(...),concat(...)))
```

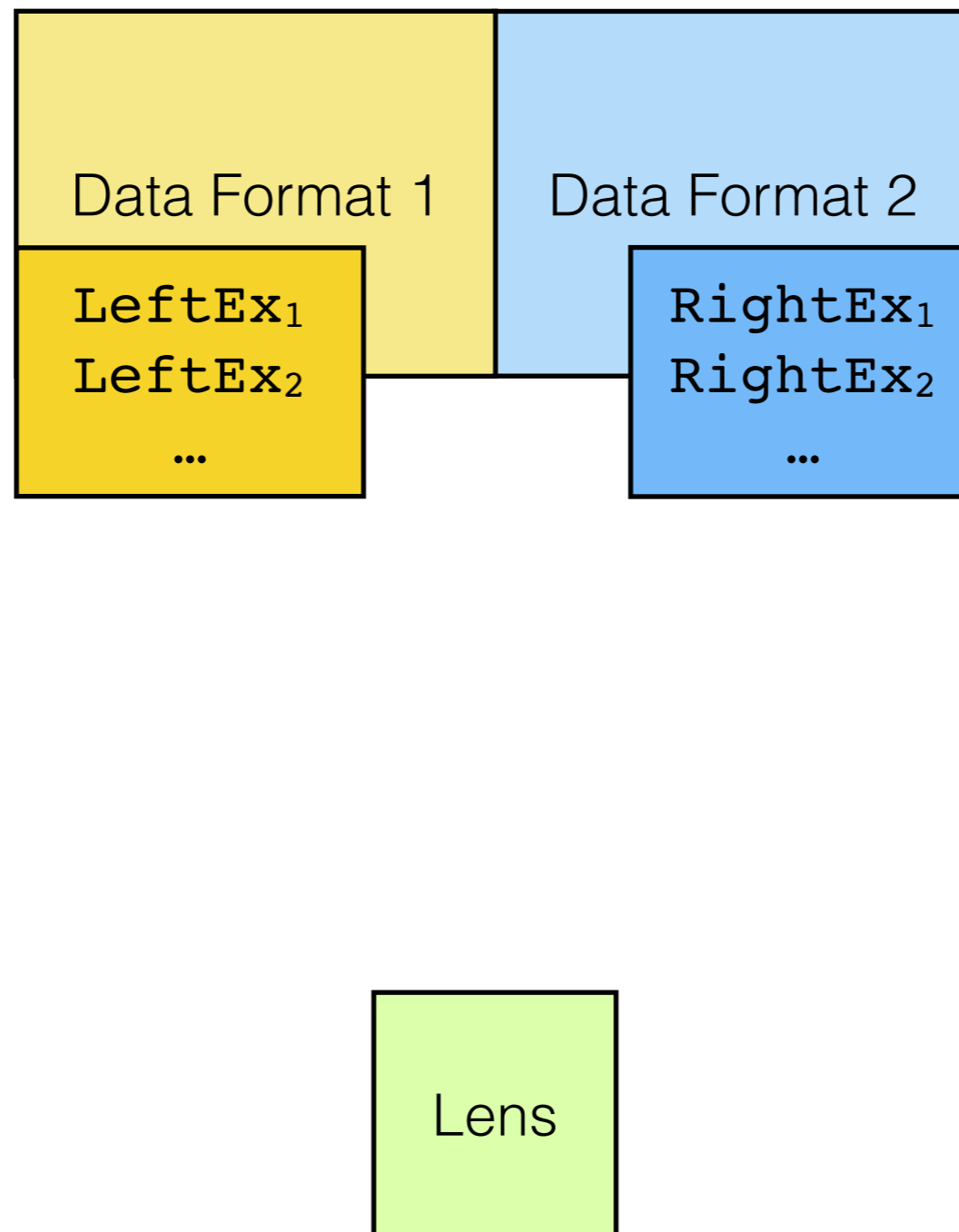
Synthesis as a Solution!



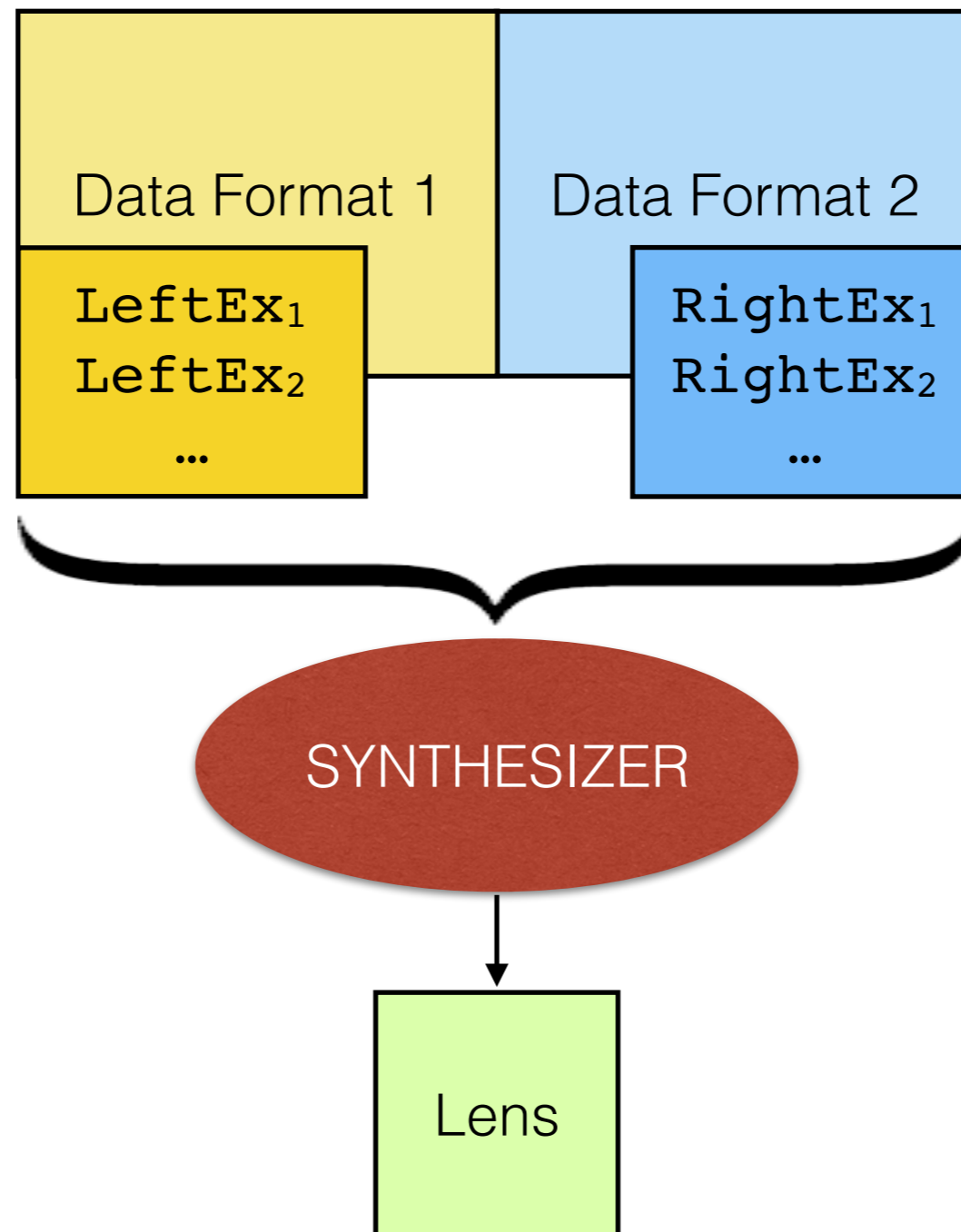
Synthesis as a Solution!



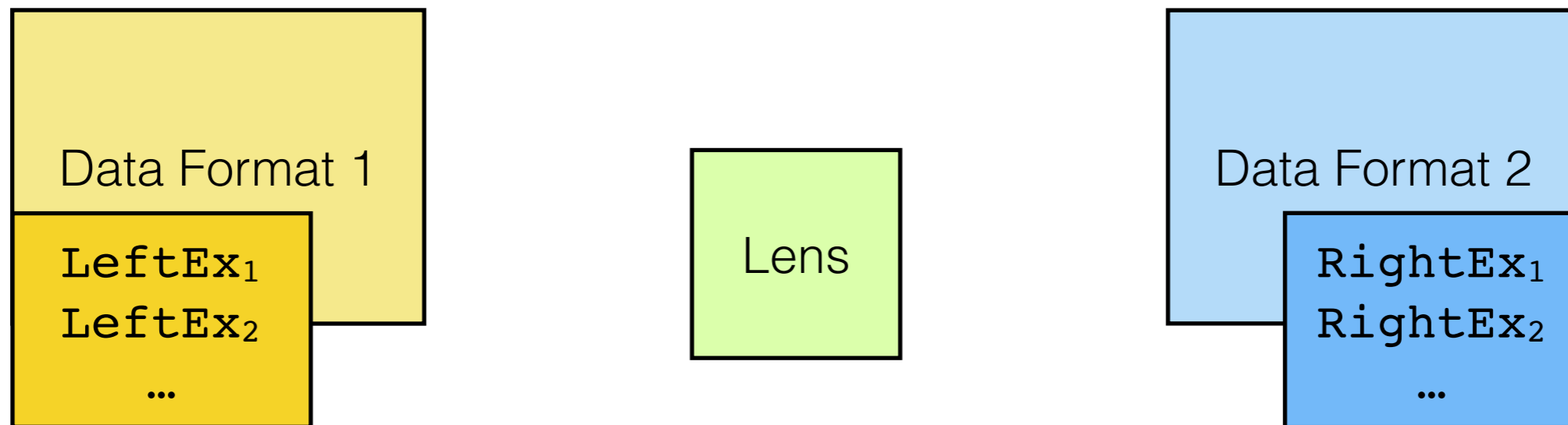
Synthesis as a Solution!



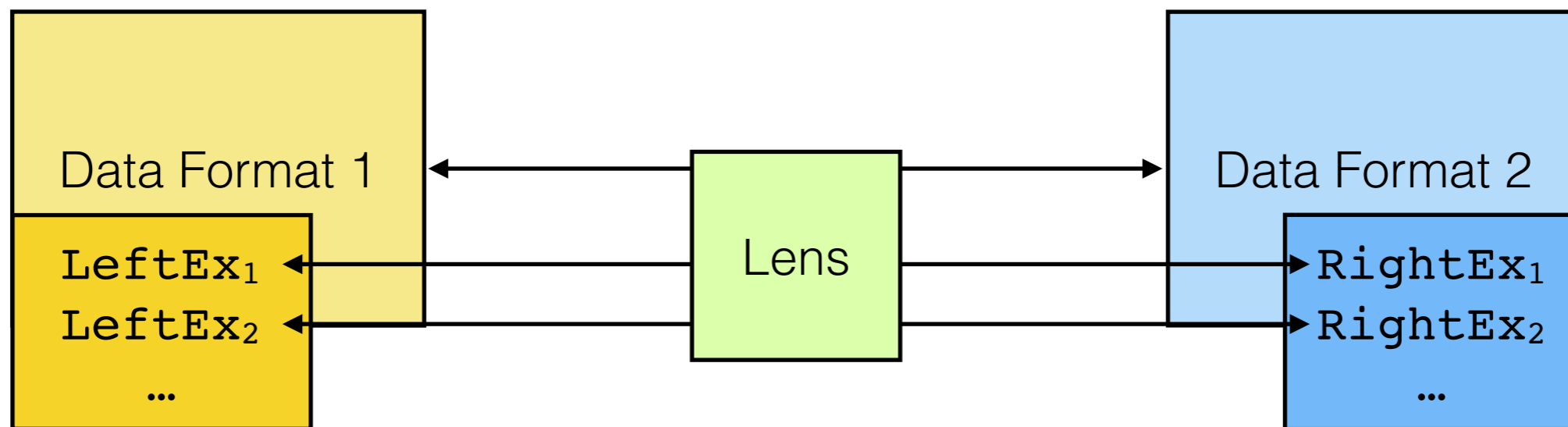
Synthesis as a Solution!



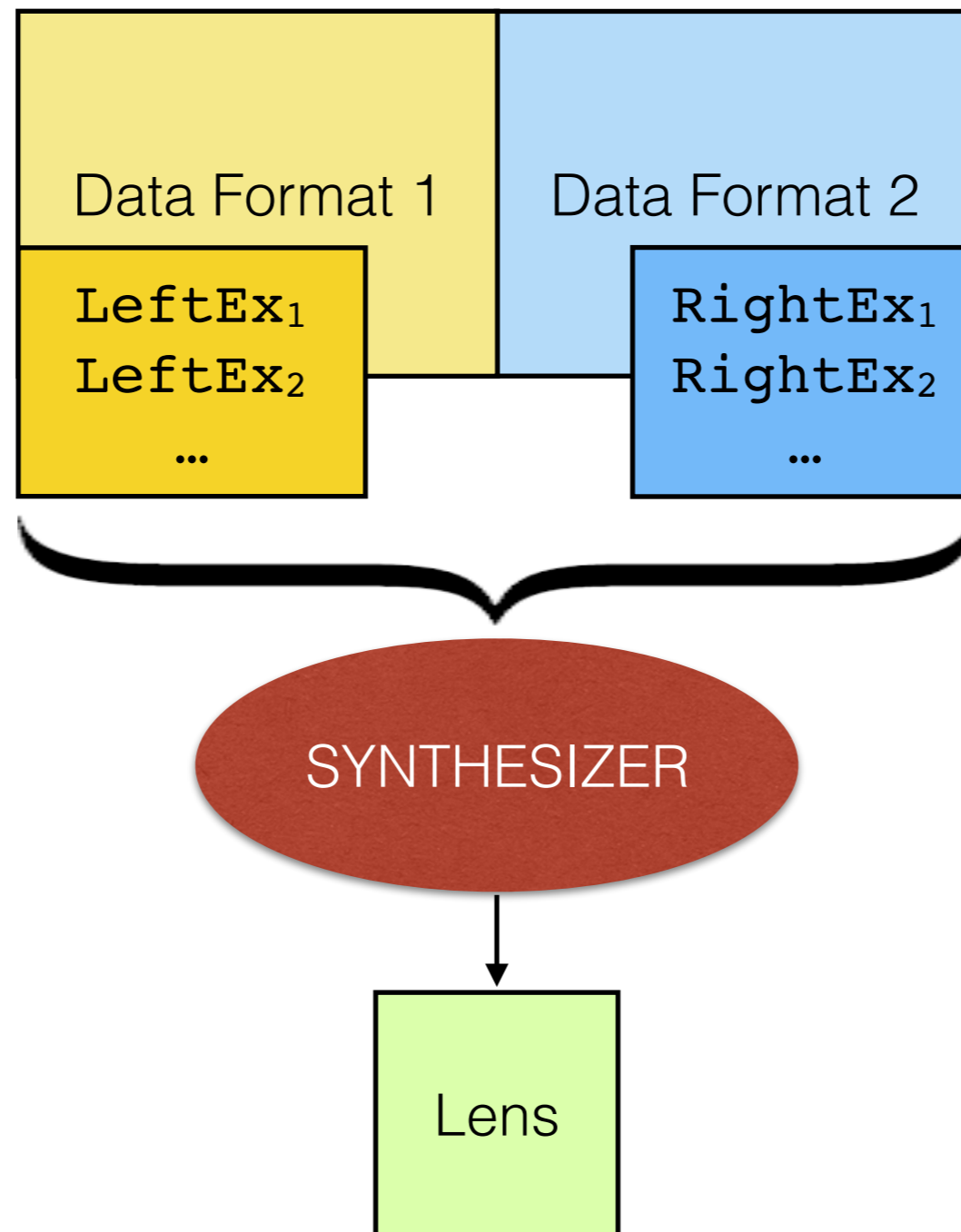
Synthesis as a Solution!



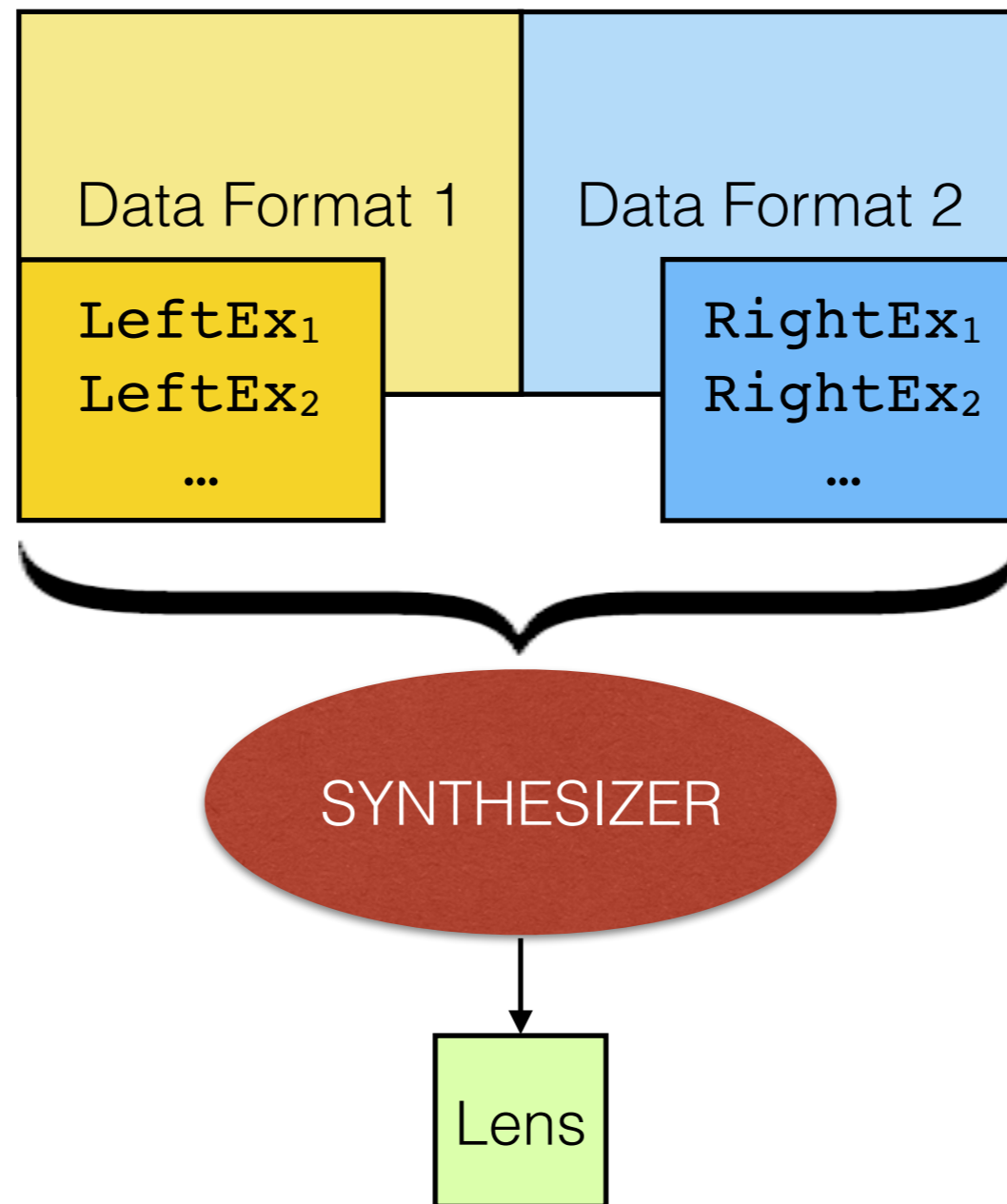
Synthesis as a Solution!



Synthesis as a Solution!



Synthesis as a Solution!



Talk Structure

Present...

- Existing lens language (bijective boomerang)
- Naive approach to synthesis, to motivate our approach
- **DNF lenses** which are more suited to synthesis
- **Algorithm** for synthesizing DNF lenses from types and examples
- **Evaluation of our algorithm** of a test suite of 25 synthesis tasks

Bijjective Boomerang

Bijjective Boomerang

- $l : R \Leftrightarrow S$
- $l.get : R \rightarrow S$
- $l.put : S \rightarrow R$
- Invertibility Guarantees?
 - $l.put \circ l.get = id|_R$
 - $l.get \circ l.put = id|_S$
 - From side conditions
- Extended example with BibTex to EndNote author field

Constant Lens

```
author={Kleene, Stephen Cole}
```

```
%A Stephen Cole Kleene
```


Constant Lens

```
author={Kleene, Stephen Cole}
```

```
%A Stephen Cole Kleene
```

```
const ( "author={", "%A" ) : "author=" <=> "%A"
```

Constant Lens

`author={Kleene, Stephen Cole}`

`%A Stephen Cole Kleene`

`const("author={", "%A") : "author=" <=> "%A"`

`const(s, t) : s <=> t`

Identity Lens

```
author={Kleene, Stephen Cole}
```

```
%A Stephen Cole Kleene
```

Identity Lens

```
author={Kleene, Stephen Cole}
```

```
%A Stephen Cole Kleene
```

```
id([A-Z][a-z]*) : [A-Z][a-z]* <=> [A-Z][a-z]*
```

Identity Lens

author={Kleene, Stephen Cole}

%A Stephen Cole Kleene

$\text{id}([A-Z][a-z]^*) : [A-Z][a-z]^* \Leftrightarrow [A-Z][a-z]^*$

side conditions

$\text{id}(R) : R \Leftrightarrow R$

Concat Lens

```
author={Kleene, Stephen Cole}
```

```
%A Stephen Cole Kleene
```

Concat Lens

`author={Kleene, Stephen Cole}`

`%A Stephen Cole Kleene`

`concat(l1, l2) : R1.R2 <=> S1.S2`

Concat Lens

author={Kleene, Stephen Cole}

%A Stephen Cole Kleene

concat(l_1, l_2) : $R_1 \cdot R_2 \Leftrightarrow S_1 \cdot S_2$

l_1 : $R_1 \Leftrightarrow S_1$

l_2 : $R_2 \Leftrightarrow S_2$

side conditions

concat(l_1, l_2) : $R_1 \cdot R_2 \Leftrightarrow S_1 \cdot S_2$

Swap Lens

```
author={Kleene, Stephen Cole}
```

```
%A Stephen Cole Kleene
```

Swap Lens

author={Kleene, Stephen Cole}

%A Stephen Cole Kleene

swap(l_1, l_2) : $R_1 \cdot R_2 \Leftrightarrow S_2 \cdot S_1$

Swap Lens

author={Kleene, Stephen Cole}

%A Stephen Cole Kleene

swap(l_1, l_2) : $R_1 \cdot R_2 \Leftrightarrow S_2 \cdot S_1$

l_1 : $R_1 \Leftrightarrow S_1$

l_2 : $R_2 \Leftrightarrow S_2$

side conditions

swap(l_1, l_2) : $R_1 \cdot R_2 \Leftrightarrow S_2 \cdot S_1$

Or Lens

ε
author={Kleene, Stephen Cole}

ε
%A Stephen Cole Kleene

Or Lens

ε
author={Kleene, Stephen Cole}

ε
%A Stephen Cole Kleene

or(l_1, l_2) : $R_1 | R_2 \Leftrightarrow S_1 | S_2$

Or Lens

ε
author={Kleene, Stephen Cole}

ε
%A Stephen Cole Kleene

or(l_1, l_2) : $R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2$

l_1 : $R_1 \Leftrightarrow S_1$

l_2 : $R_2 \Leftrightarrow S_2$

side conditions

or(l_1, l_2) : $R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2$

Iterate Lens

```
author={Hofmann, Martin  
and Pierce, Benjamin  
and Wagner, Daniel}
```

```
%A Martin Hofmann  
%A Benjamin Pierce  
%A Daniel Wagner
```

Iterate Lens

```
author={Hofmann, Martin  
and Pierce, Benjamin  
and Wagner, Daniel}
```

```
%A Martin Hofmann  
%A Benjamin Pierce  
%A Daniel Wagner
```

```
iterate(l) : R* <=> S*
```


Iterate Lens

```
author={Hofmann, Martin  
and Pierce, Benjamin  
and Wagner, Daniel}
```

```
%A Martin Hofmann  
%A Benjamin Pierce  
%A Daniel Wagner
```

```
iterate(l) : R* <=> S*
```

$$l : R \Leftrightarrow S$$

side conditions

```
iterate(l) : R* <=> S*
```

Single Author Lens!

```
author={Kleene, Stephen Cole},
```

```
%A Stephen Cole Kleene
```

Single Author Lens!

```
author={Kleene, Stephen Cole},           %A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"},  
  1)
```

Single Author Lens!

```
author={Kleene, Stephen Cole},
```

```
%A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"),  
concat(  
  1,  
  const("},",  $\epsilon$ )))
```

Single Author Lens!

```
author={Kleene, Stephen Cole},           %A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"),  
concat(  
  1,  
  const("}", "", ε)))  
id(" "[A-Z][a-z]*)
```

Single Author Lens!

```
author={Kleene, Stephen Cole},           %A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"),  
concat(  
  1,  
  const("}", " , ε"))
```

```
iterate(id(" "[A-Z][a-z]*))
```

Single Author Lens!

```
author={Kleene, Stephen Cole},           %A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"),  
concat(  
  swap(  
    1,  
    iterate(id(" "[A-Z][a-z]*))),  
  const("},", "\u03b5"))
```

Single Author Lens!

```
author={Kleene, Stephen Cole},           %A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"),  
concat(  
  swap(  
    concat(  
      id("[A-Z][a-z]*"),  
      1),  
    iterate(id(" "[A-Z][a-z]*))),  
  const("}", " , ε"))
```


Single Author Lens!

```
author={Kleene, Stephen Cole},           %A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"),  
concat(  
  swap(  
    concat(  
      id("[A-Z][a-z]*"),  
      const(", ",  $\epsilon$ )),  
      iterate(id(" "[A-Z][a-z]*))),  
    const("}", ",  $\epsilon$ ))
```

Single Author Lens!

```
author={Kleene, Stephen Cole},
```

```
%A Stephen Cole Kleene
```

```
concat(  
  const("author={", "%A"),  
  concat(  
    swap(  
      concat(  
        id("[A-Z][a-z]*"),  
        const(", ",  $\epsilon$ ),  
        iterate(id(" "[A-Z][a-z]*)),  
        const("}", ",  $\epsilon$ ))  
    )  
  )  
)
```

-
-

```
"author={"  
  [A-Z][a-z]  
  ", "  
  , "  
  (" "[A-Z][a-z])*  
  }"  
  
<=>  
  
  "%A"  
  (" "[A-Z][a-z])*  
  [A-Z][a-z]  
   $\epsilon$   
   $\epsilon$ 
```

Single Author Lens!

`author={Kleene, Stephen Cole},` `%A Stephen Cole Kleene`

```
concat(  
  const("author={", "%A"),  
  concat(  
    swap(  
      concat(  
        id("[A-Z][a-z]*"),  
        const(", ",  $\epsilon$ ),  
        iterate(id(" "[A-Z][a-z]*)),  
        const("}", ",  $\epsilon$ ))  
    )  
  )  
)
```

```
"author={"  
  ((([A-Z][a-z]  
    ",")  
    (" "[A-Z][a-z])*  
  "}")  
)
```

●
● \Leftrightarrow

```
"%A"  
((((" "[A-Z][a-z])*  
  ([A-Z][a-z]  
     $\epsilon$ ))  
   $\epsilon$ )  
)
```

Potential Problem!

```
"author={"  
  [A-Z][a-z]  
  ", "  
  (" "[A-Z][a-z])*  
  "}"
```

<=>

```
"%A"  
  (" "[A-Z][a-z])*  
  [A-Z][a-z]  
  ε  
  ε
```

```
"author={"  
  ((([A-Z][a-z]  
    ", ")  
    (" "[A-Z][a-z])*)  
  "}"
```

<=>

```
"%A"  
  ((((" "[A-Z][a-z])*  
    ([A-Z][a-z]  
    ε))  
  ε)
```

Potential Problem!

```
"author="{  
  ([A-Z][a-z]  
  (" , "  
    (" "[A-Z][a-z])  
    "}")
```

<=>

```
"%A"  
  (" "[A-Z][a-z])  
  ([A-Z][a-z]  
  (ε  
  ε)))
```

```
"author="{  
  ((([A-Z][a-z]  
    " , "  
    (" "[A-Z][a-z])  
    "}")
```

<=>

```
"%A"  
  ((((" "[A-Z][a-z])  
    ([A-Z][a-z]  
    ε))  
  ε)
```

Potential Problem!

```
"author={"  
  [A-Z][a-z]  
  "," "  
  (" "[A-Z][a-z])*  
  "}"
```

<=>

```
"%A"  
(" "[A-Z][a-z])*  
[A-Z][a-z]  
ε  
ε
```

```
"author={"  
  ((([A-Z][a-z]  
    "," ")  
    (" "[A-Z][a-z])*)  
  "}"
```

<=>

```
"%A"  
((((" "[A-Z][a-z])*  
  ([A-Z][a-z]  
    ε))  
  ε)
```

Potential Problem!

Name = [A-Z][a-z]

```
"author={"  
  Name  
  ","  
  (" "Name)*  
  "}"
```

<=>

```
"%A"  
(" "Name)*  
  Name  
  ε  
  ε
```

```
"author={"  
  ((([A-Z][a-z]  
    ",")  
    (" "[A-Z][a-z])* )  
  "}"
```

<=>

```
"%A"  
((((" "[A-Z][a-z])*  
  ([A-Z][a-z]  
    ε))  
  ε)
```

Potential Problem!

```
Name = [A-Z][a-z]  
Names = (" "Name)*
```

```
"author={ "  
  Name  
  " , "  
  Names  
  " } "
```

<=>

```
"%A"  
  Names  
  Name  
  ε  
  ε
```

```
"author={ "  
  ((([A-Z][a-z]  
    " , "  
    (" "[A-Z][a-z])* )  
  " } "
```

<=>

```
"%A"  
  ((((" "[A-Z][a-z])*  
    ([A-Z][a-z]  
    ε))  
  ε)
```


Potential Problem!

```
Name = [A-Z][a-z]
Names = (" "Name)*
```

```
SingleBibAuth =
  "author={" Name "," Names "}"
```

```
SingleEndAuth =
  "%A" Names Name
```

```
SingleBibAuth <=> SingleEndAuth
```

```
"author={"
  ((([A-Z][a-z]
    ","))
    (" "[A-Z][a-z])*
  "}")
<=>
"%A"
((((" "[A-Z][a-z])*
  ([A-Z][a-z]
    ε))
  ε)
```

Retrying

```
Name = [A-Z][a-z]
Names = (" "Name)*
```

```
SingleBibAuth =
  "author={" Name "," Names "}"
```

```
SingleEndAuth =
  "%A" Names Name
```

```
SingleBibAuth <=> SingleEndAuth
```

```
"author={"
  ((([A-Z][a-z]
    ","))
  (" "[A-Z][a-z])*
  "}")
<=>
"%A"
(((("[A-Z][a-z])*
  ([A-Z][a-z]
  ε))
  ε)
```

$$R \equiv R'$$

$$S \equiv S'$$

$$\frac{l : R' \Leftrightarrow S'}{l : R \Leftrightarrow S}$$

$$\mathbf{R} \equiv \mathbf{R}'$$

- Semiring Axioms

$$R \equiv R'$$

- Semiring Axioms

- $R \mid \emptyset \equiv R$

- $R \cdot \emptyset \equiv \emptyset$

- $\emptyset \cdot R \equiv \emptyset$

- $(R \cdot R') \cdot R'' \equiv R \cdot (R' \cdot R'')$

- $(R \mid R') \mid R'' \equiv R \mid (R' \mid R'')$

- $R \mid T \equiv T \mid R$

- $R \cdot (R' \mid R'') \equiv (R \cdot R') \mid (R \cdot R'')$

- $(R' \mid R'') \cdot R \equiv (R' \cdot R) \mid (R'' \cdot R)$

- $\varepsilon \cdot R \equiv R$

- $R \cdot \varepsilon \equiv R$

$$R \equiv R'$$

- Semiring Axioms
- ε -unrolling for $*$

$$R \equiv R'$$

- Semiring Axioms
- ε -unrolling for $*$
 - $R^* \equiv \varepsilon \mid (R \cdot R^*)$
 - $R^* \equiv \varepsilon \mid (R^* \cdot R)$

$$\mathbf{R} \equiv \mathbf{R}'$$

- Semiring Axioms
- ε -unrolling for $*$
- \mathbf{R} is *definitionally equivalent* to \mathbf{R}'

$$R \equiv R'$$

- Semiring Axioms
- ε -unrolling for $*$
- R is *definitionally equivalent* to R'
- Not full equivalence, but sufficient in practice!

Composition

author={Kleene, Stephen Cole}

↕ l_1

%A Stephen Cole Kleene

↕ l_2

%A1 Stephen Cole Kleene

Composition

author={Kleene, Stephen Cole}

↕ l_1

%A Stephen Cole Kleene

↕ l_2

%A1 Stephen Cole Kleene

$l_1; l_2$: SingleBibAuth \Leftrightarrow NewSingleEndAuth

Composition

author={Kleene, Stephen Cole}

$\updownarrow l_1$

%A Stephen Cole Kleene

$\updownarrow l_2$

%A1 Stephen Cole Kleene

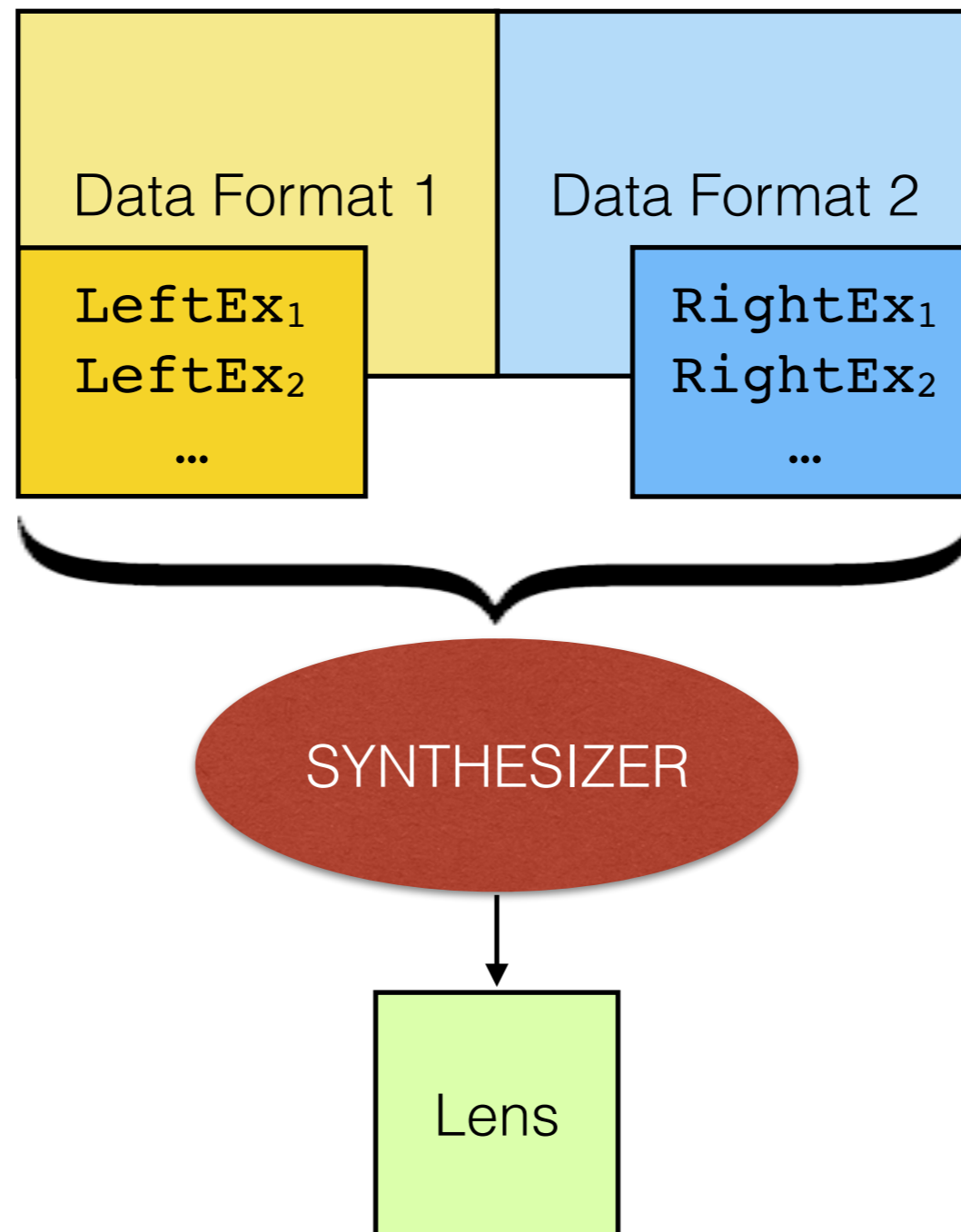
$l_1; l_2 : \text{SingleBibAuth} \Leftrightarrow \text{NewSingleEndAuth}$

$l_1 : R \Leftrightarrow S$

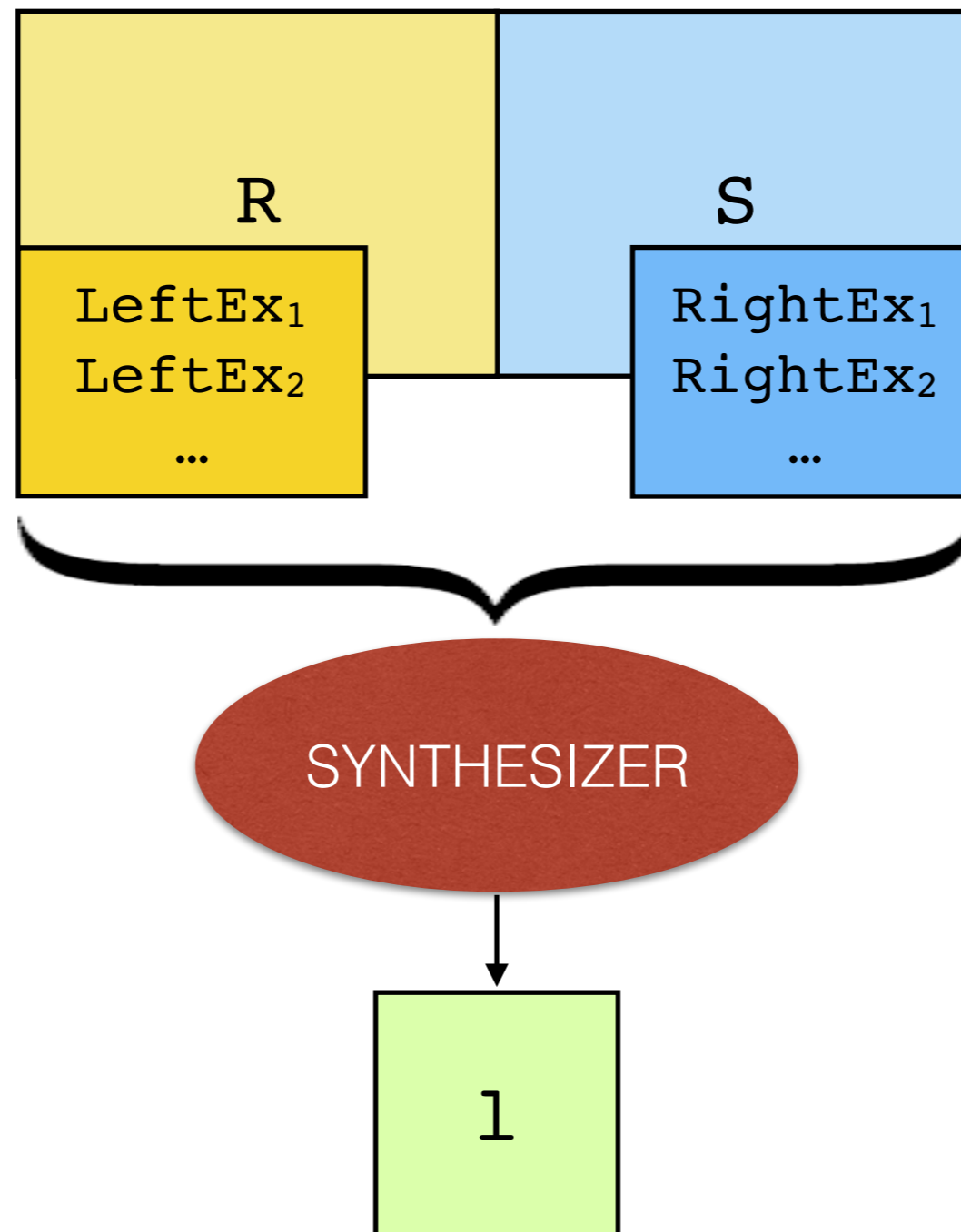
$l_2 : S \Leftrightarrow T$

$l_1; l_2 : R \Leftrightarrow T$

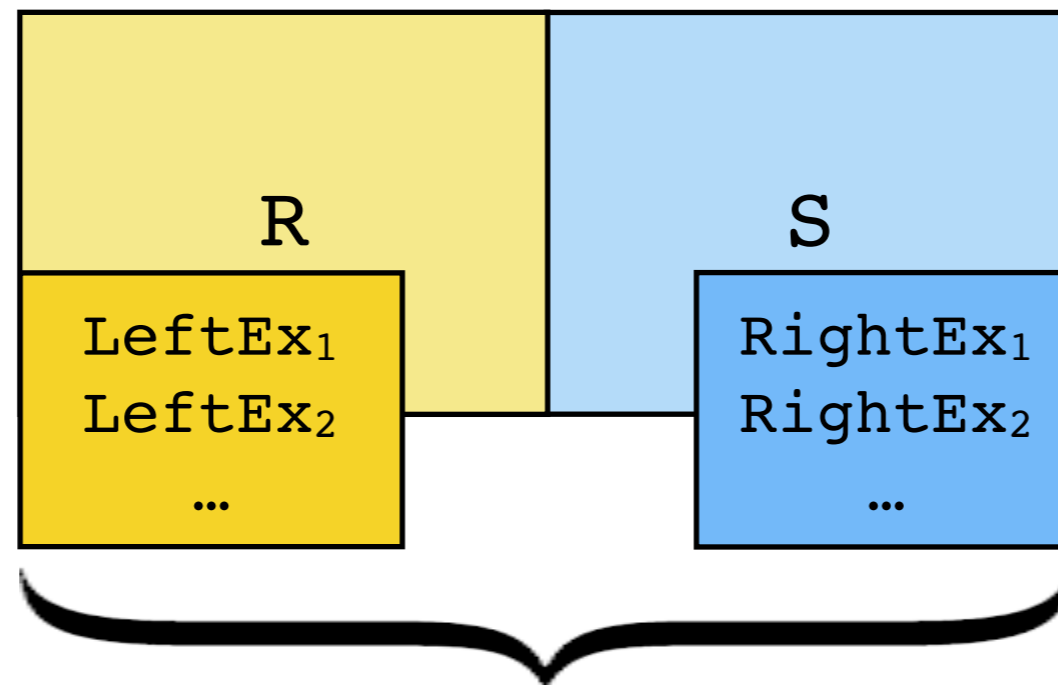
Lenses from Data Formats



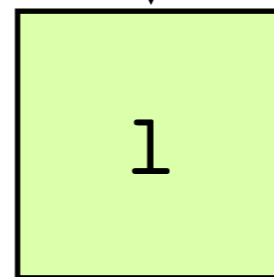
Lenses from Data Formats



Lenses from Data Formats

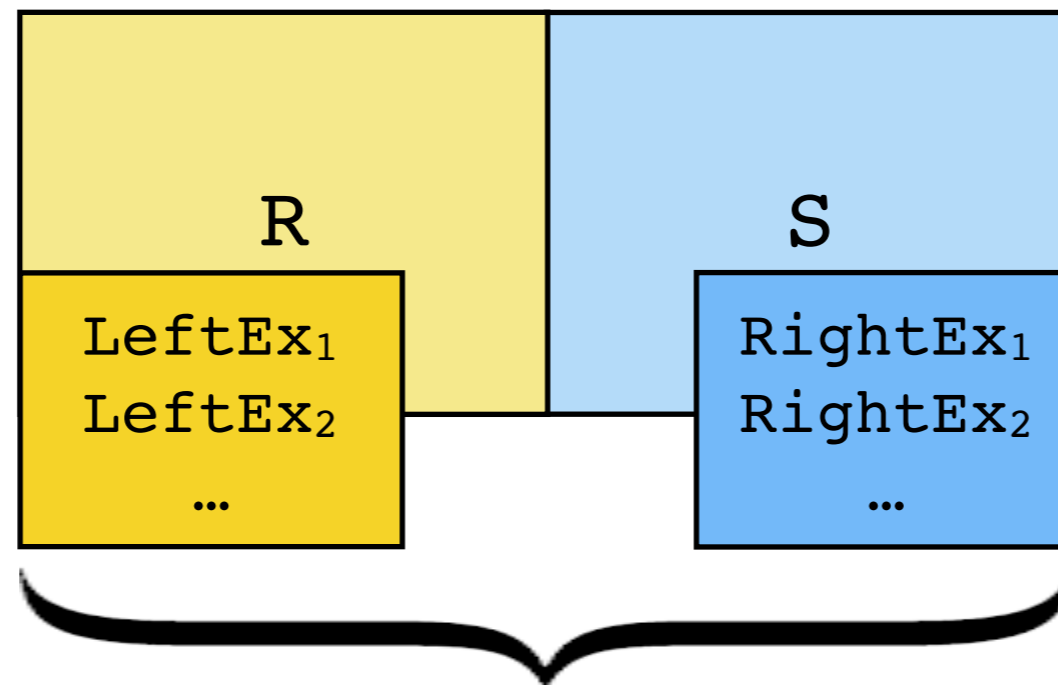


SYNTHESIZER



- $l : R \Leftrightarrow S$
- $l.get \text{ LeftEx}_i = \text{RightEx}_i$

Lenses from Data Formats



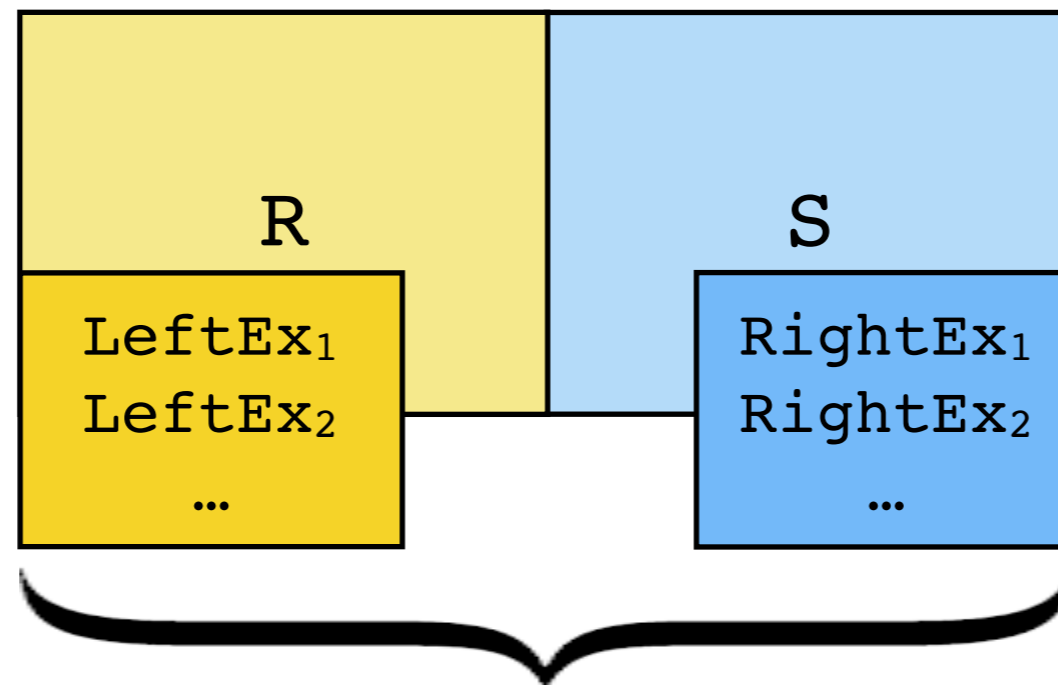
Use the type syntax to direct the term creation!

SYNTHESIZER

l

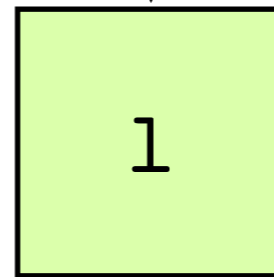
- $l : R \Leftrightarrow S$
- $l.get\ LeftEx_i = RightEx_i$

Lenses from Data Formats



Use the type syntax to direct the term creation!

? : R <=> S



- $l : R \lt;=> S$
- $l.get\ LeftEx_i = RightEx_i$

Naive Algorithm

The Easy Cases!

$$? : R_1 \mid R_2 \iff S_1 \mid S_2$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$? : R^* \Leftrightarrow S^*$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$
$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$
$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$? : R_1 \cdot R_2 \Leftrightarrow S_1 \cdot S_2$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1 \cdot R_2 \Leftrightarrow S_1 \cdot S_2}$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$? : R_1.R_2 \Leftrightarrow S_1.S_2$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$? : s \Leftrightarrow t$$

The Easy Cases!

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\text{const}(s, t) : s \Leftrightarrow t$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$? : R_1 \cdot (R_2 \mid R_3) \Leftrightarrow (R_1 \cdot R_2) \mid (R_1 \cdot R_3)$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{? : (R_1 \cdot R_2) \mid (R_1 \cdot R_3) \Leftrightarrow (R_1 \cdot R_2) \mid (R_1 \cdot R_3)}{? : R_1 \cdot (R_2 \mid R_3) \Leftrightarrow (R_1 \cdot R_2) \mid (R_1 \cdot R_3)}$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$? : R_1 \cdot (R_2 \cdot R_3) \Leftrightarrow (R_1 \cdot R_2) \cdot R_3$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : R_1 \cdot (R_2 \cdot R_3) \Leftrightarrow R_1 \cdot (R_2 \cdot R_3)}{? : R_1 \cdot (R_2 \cdot R_3) \Leftrightarrow (R_1 \cdot R_2) \cdot R_3}$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{? : ?_1 \Leftrightarrow ?_2 \quad ?_1 \equiv R \quad ?_2 \equiv S}{? : R \Leftrightarrow S}$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : ?_1 \Leftrightarrow ?_2 \quad ?_1 \equiv R \quad ?_2 \equiv S}{? : R \Leftrightarrow S}$$

$$? : R1 \cdot R2 \cdot R3 \cdot R4 \Leftrightarrow R3 \cdot R1 \cdot R4 \cdot R2$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : ?_1 \Leftrightarrow ?_2 \quad ?_1 \equiv R \quad ?_2 \equiv S}{? : R \Leftrightarrow S}$$

$$\frac{?_1 : R \Leftrightarrow ? \quad ?_2 : ? \Leftrightarrow S}{?_1 ; ?_2 : R \Leftrightarrow S}$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : ?_1 \Leftrightarrow ?_2 \quad ?_1 \equiv R \quad ?_2 \equiv S}{? : R \Leftrightarrow S}$$

$$\frac{?_1 : R \Leftrightarrow ? \quad ?_2 : ? \Leftrightarrow S}{?_1; ?_2 : R \Leftrightarrow S}$$

The Hard Cases

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{or}(?_1, ?_2) : R_1 \mid R_2 \Leftrightarrow S_1 \mid S_2}$$

$$\frac{? : R \Leftrightarrow S}{\text{iterate}(?) : R^* \Leftrightarrow S^*}$$

$$\frac{}{\text{const}(s, t) : s \Leftrightarrow t}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_1 \quad ?_2 : R_2 \Leftrightarrow S_2}{\text{concat}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{?_1 : R_1 \Leftrightarrow S_2 \quad ?_2 : R_2 \Leftrightarrow S_1}{\text{swap}(?_1, ?_2) : R_1.R_2 \Leftrightarrow S_1.S_2}$$

$$\frac{? : ?_1 \Leftrightarrow ?_2 \quad ?_1 \equiv R \quad ?_2 \equiv S}{? : R \Leftrightarrow S}$$

$$\frac{?_1 : R \Leftrightarrow ? \quad ?_2 : ? \Leftrightarrow S}{?_1; ?_2 : R \Leftrightarrow S}$$

HUGE SEARCH SPACE

Approach

- Reduce Equivalences
 - Alternate language of types: *DNF Regular Expressions*
 - Restrict where retyping can be used
- Remove Composition
 - Alternate language of terms: *DNF Lenses*
 - DNF lenses closed under composition

DNF Regular Expressions and DNF Lenses

DNF Regular Expressions

- Fully distributed regular expression
- No associativity information
- Base strings enforced at fixed locations

Converting to DNF Regular Expressions ↓

"a" ("b" | ("c" "d")*) "e"* "f"

Converting to DNF Regular Expressions ↓

"a" ("b" | ("c" ↓ "d")*) "e"* "f"
"a" ("b" | "cd"*) "e"* "f"

Converting to DNF Regular Expressions ↓

"a" ("b" | ("c" "d")*) "e"* "f"

"a" ("b" | "cd"*) "e"* "f"

(("a" "b") | ("a" "cd"*)) "e"* "f"

Converting to DNF Regular Expressions ↓

"a" ("b" | ("c" "d")*) "e"* "f"
↓
"a" ("b" | "cd"*) "e"* "f"
↓
(("a" "b") | ("a" "cd"*)) "e"* "f"
↓
("ab" | ("a" "cd"*)) "e"* "f"

Converting to DNF Regular Expressions ↓

"a" ("b" | ("c" "d")*) "e"* "f"
↓
"a" ("b" | "cd"*) "e"* "f"
↓
(("a" "b") | ("a" "cd"*)) "e"* "f"
↓
("ab" | ("a" "cd"*)) "e"* "f"
↓
("ab" "ef") | ("a" "cd"* "e"* "f")

Converting to DNF Regular Expressions ↓

"a" ("b" | ("c" "d")*) "e"* "f"
"a" ("b" | "cd"*) "e"* "f"
(("a" "b") | ("a" "cd"*)) "e"* "f"
("ab" | ("a" "cd"*)) "e"* "f"
("ab" "ef") | ("a" "cd"* "e"* "f")
"abef" | ("a" "cd"* "e"* "f")

Converting to DNF Regular Expressions \Downarrow

"abef" | ("a" "cd"* "e"* "f")

Atoms

$A ::= DR^*$

iterate

Sequences

$SQ ::= \langle s_0; A_1; \dots; A_n; s_n \rangle$

sequence

DNF Regular Expressions

$DR ::= \{SQ_1; \dots; SQ_n\}$

choose

Converting to DNF Regular Expressions \Downarrow

"abef" | ("a" "cd"* "e"* "f")
 \Downarrow
{ "abef" ; "a" "cd"* "e"* "f" }

Atoms

$A ::= DR^*$

iterate

Sequences

$SQ ::= \langle s_0 ; A_1 ; \dots ; A_n ; s_n \rangle$

sequence

DNF Regular Expressions

$DR ::= \{ SQ_1 ; \dots ; SQ_n \}$

choose

Converting to DNF Regular Expressions \Downarrow

"abef" | ("a" "cd"* "e"* "f")
↓
{ "abef" ; "a" "cd"* "e"* "f" }
↓
{ <"abef"> ; <"a"; "cd"*; "e"*; "f"> }

Atoms

$A ::= DR^*$

iterate

Sequences

$SQ ::= \langle s_0; A_1; \dots; A_n; s_n \rangle$

sequence

DNF Regular Expressions

$DR ::= \{ SQ_1; \dots; SQ_n \}$

choose

Converting to DNF Regular Expressions \Downarrow

"abef" | ("a" "cd"* "e"* "f")
 \Downarrow
{ "abef" ; "a" "cd"* "e"* "f" }
 \Downarrow
{ <"abef"> ; <"a";"cd"*;"e"*;"f"> }
 \Downarrow
{ <"abef"> ; <"a";"cd"*;" ϵ ";"e"*;"f"> }

Atoms

$A ::= DR^*$

iterate

Sequences

$SQ ::= \langle s_0; A_1; \dots; A_n; s_n \rangle$

sequence

DNF Regular Expressions

$DR ::= \{ SQ_1; \dots; SQ_n \}$

choose

Converting to DNF Regular Expressions ↓

```
"a" ("b" | ("c" "d")*) "e"* "f"
```

```
{ <"abef"> ; <"a";"cd"*;ε;"e"*;"f"> }
```

Converting to DNF Regular Expressions ↓

```
"a" ("b" | ("c" "d")*) "e"* "f"
```

```
"a" ε ("b" | "cd"* | ∅) ("e"* "f")
```

```
{ <"abef"> ; <"a";"cd"*;ε;"e"*;"f"> }
```

Converting to DNF Regular Expressions \Downarrow

```
"a" ("b" | ("c" "d")*) "e"* "f"
```

```
"a"  $\epsilon$  ("b" | "cd"* |  $\emptyset$ ) ("e"* "f")
```

```
⋮
```

```
{ <"abef"> ; <"a";"cd"*; $\epsilon$ ;"e"*;"f"> }
```

Converting to DNF Regular Expressions \Downarrow

```
"a" ("b" | ("c" "d")*) "e"* "f"  
"a"  $\epsilon$  ("b" | "cd"* |  $\emptyset$ ) ("e"* "f")  
:  
{ <"abef"> ; <"a";"cd"*; $\epsilon$ ;"e"*;"f"> }
```

```
"abef" | ("a" ( $\epsilon$  | ("cd" "cd"*)) "e"* "f")  
:  
{ <"abef"> ; <"a";"e"*;"f"> ;  
  <"acd";"cd"*; $\epsilon$ ;"e"*;"f"> }
```

Converting to DNF Regular Expressions \Downarrow

```
"a" ("b" | ("c" "d")*) "e"* "f"  
"a"  $\epsilon$  ("b" | "cd"* |  $\emptyset$ ) ("e"* "f")  
:  
{ <"abef"> ; <"a";"cd"*; $\epsilon$ ;"e"*;"f"> }
```

```
"abef" | ("a" ( $\epsilon$  | ("cd" "cd"*)) "e"* "f")
```

:

```
{ <"abef"> ; <"a";"e"*;"f"> ;  
  <"acd";"cd"*; $\epsilon$ ;"e"*;"f"> }
```

```
("a" "cd"* "e"* "f") | "abef"
```

:

```
{ <"a";"cd"*; $\epsilon$ ;"e"*;"f"> ; <"abef"> }
```


DNF Regular Expression Rewrites

- $A | B \equiv B | A$ not part of DNF rewrites
 - DNF Lens terms address these equivalences
- $R^* \equiv \varepsilon | (R \cdot R^*)$ and $R^* \equiv \varepsilon | (R^* \cdot R)$ ARE part of DNF rewrites
 - $R^* \rightarrow \varepsilon | (R \cdot R^*)$
 - $R^* \rightarrow \varepsilon | (R^* \cdot R)$
 - $DR \rightarrow^* DR'$

DNF Lenses and Subcomponents

$dl ::= D \Leftrightarrow E$

Atom Lenses

$al ::= \text{iterate}(dl)$

iterate

Sequence Lenses

$sl ::= (\langle (s_0, t_0); al_1; \dots; al_n; (s_n, t_n) \rangle , \sigma)$

sequence

DNF Lenses

$dl ::= (\{sl_1; \dots; sl_n\} , \sigma)$

choose

Iterate Lens

```
author={Hofmann, Martin  
and Pierce, Benjamin  
and Wagner, Daniel}
```

```
%A Martin Hofmann  
%A Benjamin Pierce  
%A Daniel Wagner
```

Iterate Lens

```
author={Hofmann, Martin  
and Pierce, Benjamin  
and Wagner, Daniel}
```

```
%A Martin Hofmann  
%A Benjamin Pierce  
%A Daniel Wagner
```

```
iterate(dl) :: D* <=> E*
```

Iterate Lens

```
author={Hofmann, Martin  
and Pierce, Benjamin  
and Wagner, Daniel}
```

```
%A Martin Hofmann  
%A Benjamin Pierce  
%A Daniel Wagner
```

```
iterate(dl) :: D* <=> E*
```

$$dl :: D \Leftrightarrow E$$

side conditions

```
iterate(l) :: D* <=> E*
```

Sequence Lens

```
author={Kleene, Stephen Cole}
```

```
%A Stephen Cole Kleene
```

Sequence Lens

author={Kleene, Stephen Cole}

%A Stephen Cole Kleene

$(\langle ("author=", "%A"); a_1; (", ", \epsilon); a_2; (" } ", \epsilon) \rangle , [1;0]) ::$
 $\langle "author=" ; a_1 ; ", " ; a_2 ; " } " \rangle \Leftrightarrow \langle "%A" ; b_2 ; \epsilon ; b_1 ; \epsilon \rangle$

Sequence Lens

author={Kleene, Stephen Cole}

%A Stephen Cole Kleene

$(\langle ("author=" , "%A") ; a_{l_1} ; (" , " , \epsilon) ; a_{l_2} ; (" } " , \epsilon) \rangle , [1;0]) ::$
 $\langle "author=" ; a_1 ; " , " ; a_2 ; " } " \rangle \Leftrightarrow \langle "%A" ; b_2 ; \epsilon ; b_1 ; \epsilon \rangle$

$a_{l_1} :: a_1 \Leftrightarrow b_1$

...

$a_{l_n} :: a_n \Leftrightarrow b_n$

side conditions

$(\langle (s_0, t_0) ; a_{l_1} ; \dots ; a_{l_n} ; (s_n, t_n) \rangle , \sigma) ::$

$\langle s_0 ; a_1 ; \dots ; a_n ; s_n \rangle \Leftrightarrow \langle t_0 ; b_{\sigma_1} ; \dots ; b_{\sigma_n} ; t_n \rangle$

DNF Lens

ϵ
author={Kleene, Stephen Cole}

%A Stephen Cole Kleene
 ϵ

DNF Lens

```

                                ε
author={Kleene, Stephen Cole}    %A Stephen Cole Kleene
                                ε

```

$$(\{l_1, l_2\}, [1, 0]) :: \{R_1; R_2\} \Leftrightarrow \{S_2; S_1\}$$

DNF Lens

ε %A Stephen Cole Kleene
author={Kleene, Stephen Cole} ε

$(\{l_1, l_2\}, [1, 0]) :: \{R_1; R_2\} \Leftrightarrow \{S_2; S_1\}$

$sl_1 :: sq_1 \Leftrightarrow tq_1$

...

$sl_n :: sq_n \Leftrightarrow tq_n$

side conditions

$(\{sl_1; \dots; sl_n\}, \sigma) ::$

$\{sq_0; \dots; sq_n\} \Leftrightarrow \{tq_{\sigma 1}; \dots; tq_{\sigma n}\}$

Rewrites at End

$$\frac{\text{dl} :: D' \langle \Rightarrow \rangle E' \quad D \rightarrow^* D' \quad E \rightarrow^* E'}{\text{dl} :: D \langle \Rightarrow \rangle E}$$

Rewrites at End

$$\frac{\text{dl} :: D' \langle \Rightarrow \rangle E' \quad D \rightarrow^* D' \quad E \rightarrow^* E'}{\text{dl} :: D \langle \Rightarrow \rangle E}$$

AND NO COMPOSITION!

Equivalence Statements

Soundness

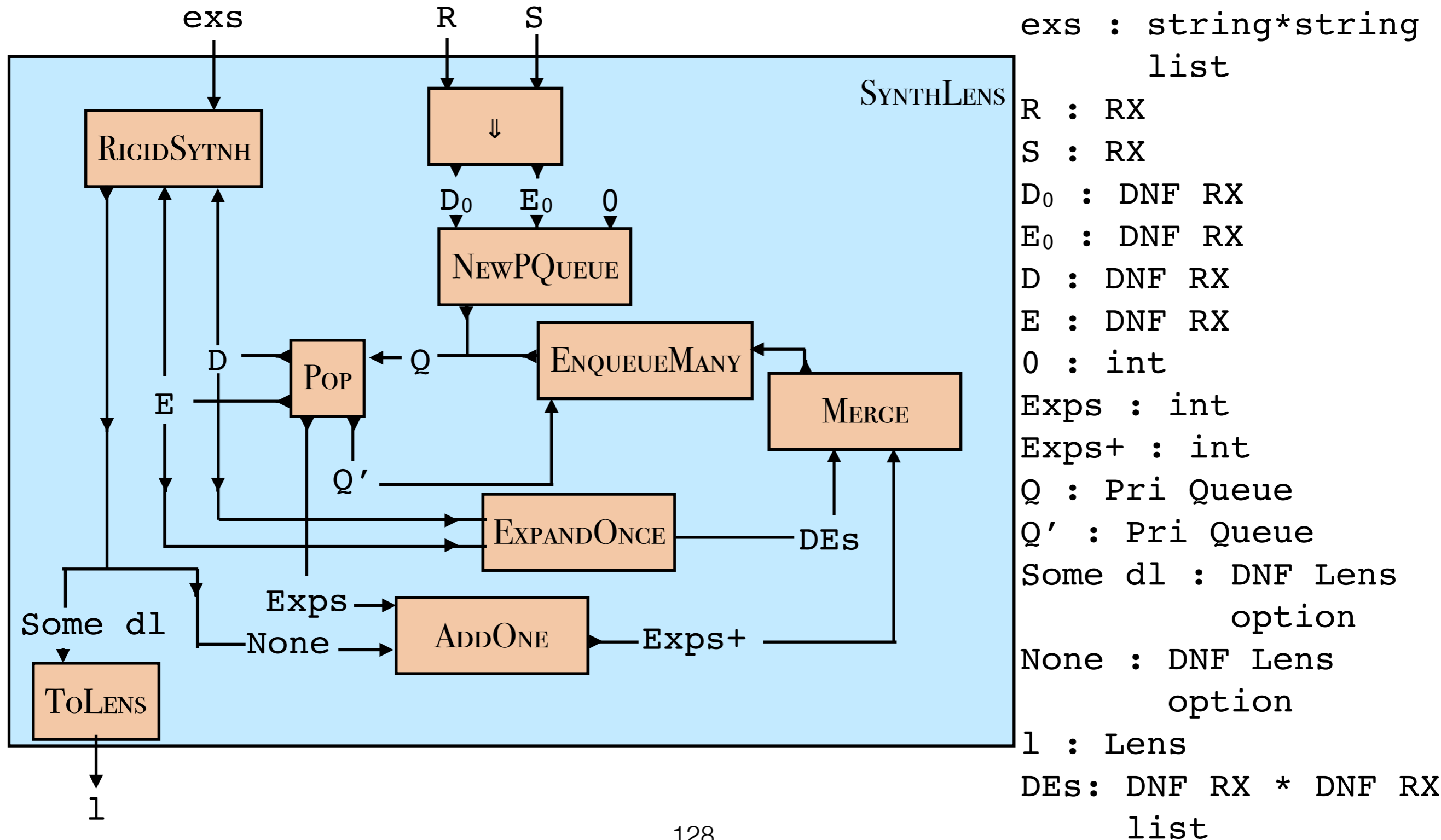
If $d1 :: DR \langle == \rangle DS$ then there exist l, R, S such that
 $l : R \langle == \rangle S$, with the same semantics

Completeness

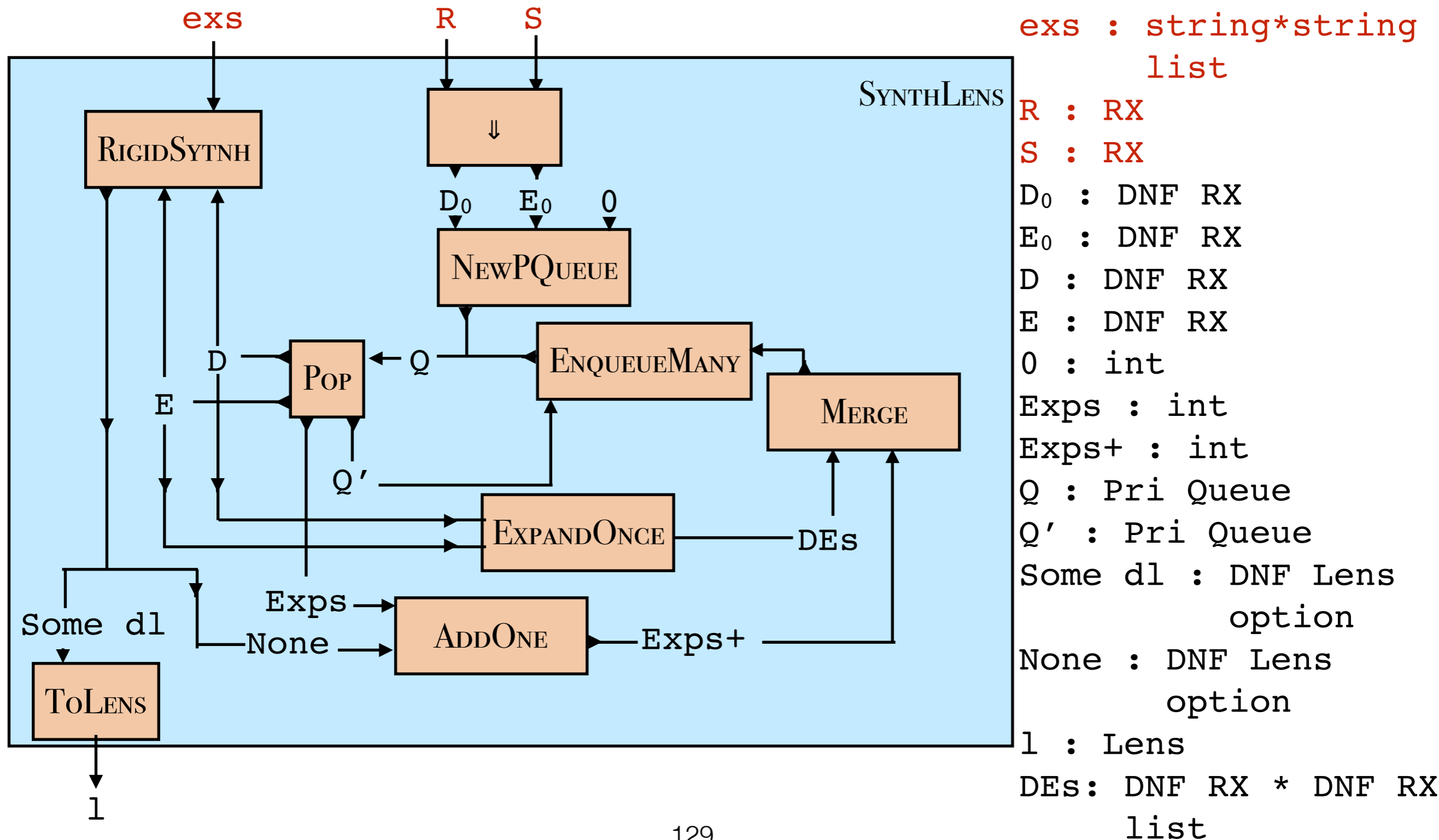
If $l : R \langle == \rangle S$ then there exists $d1$ such that
 $d1 :: \Downarrow(R) \langle == \rangle \Downarrow(S)$, with the same semantics

Synthesis Algorithm

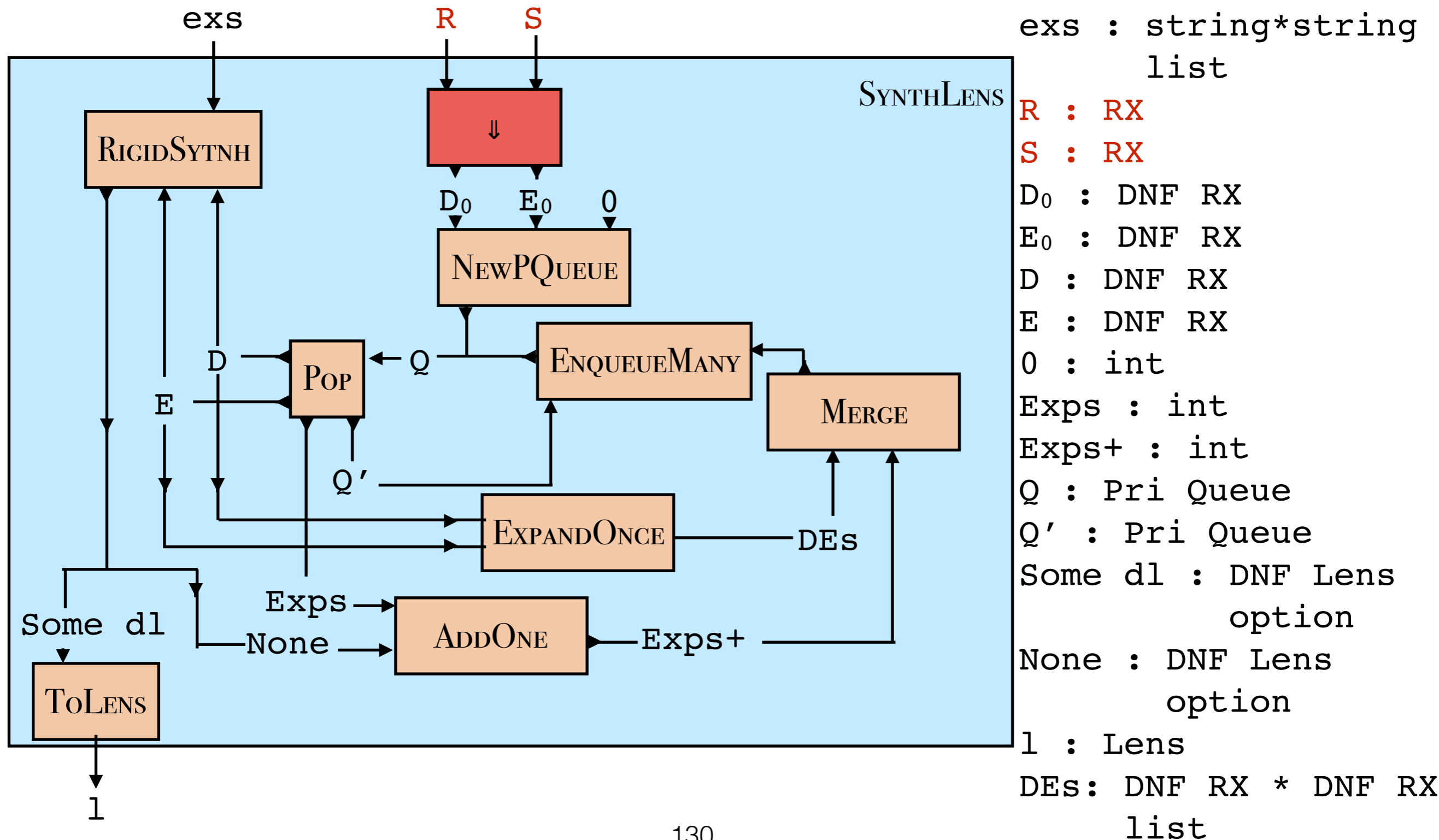
Synthesis Diagram



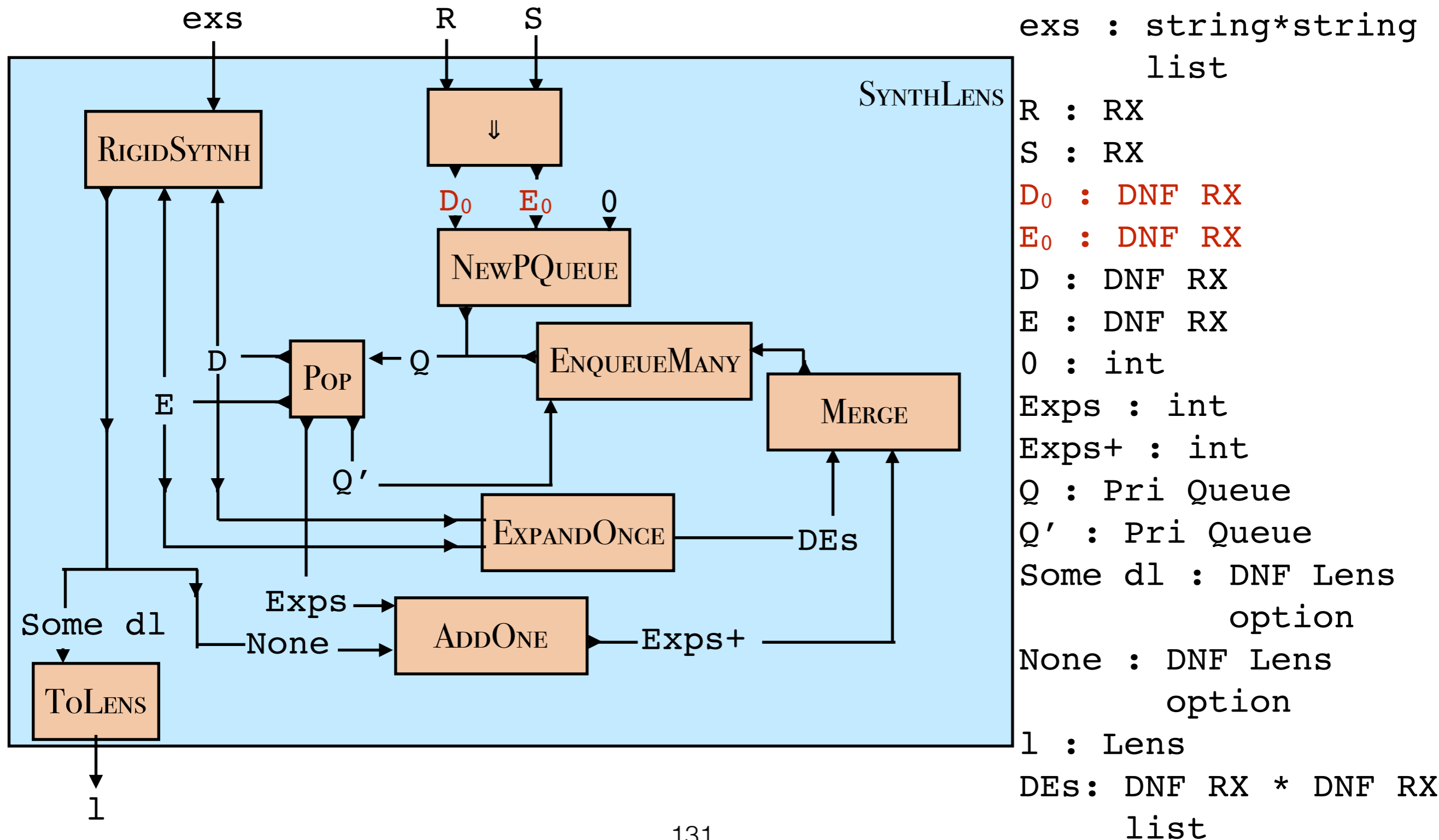
Synthesis Diagram



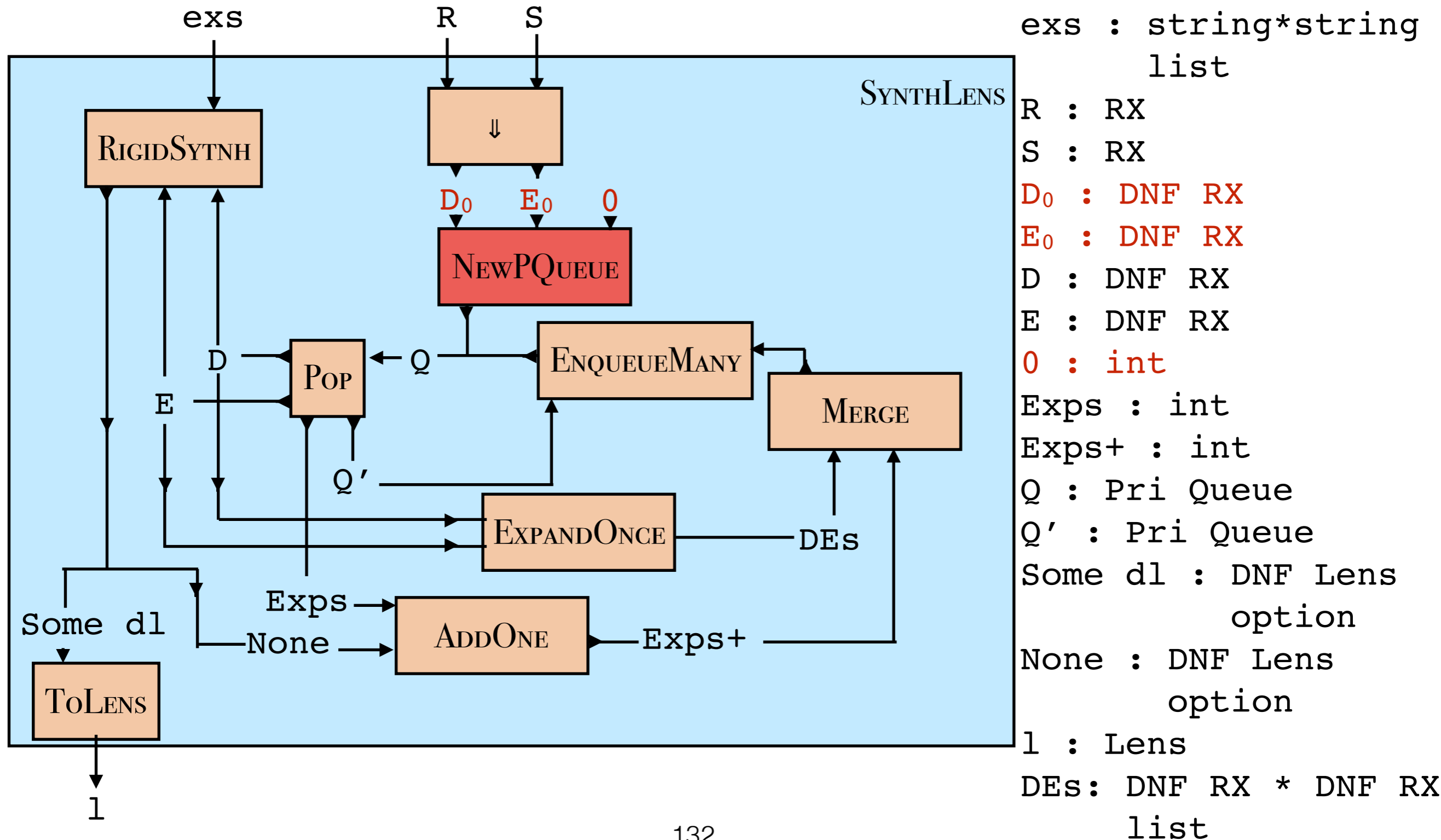
Synthesis Diagram



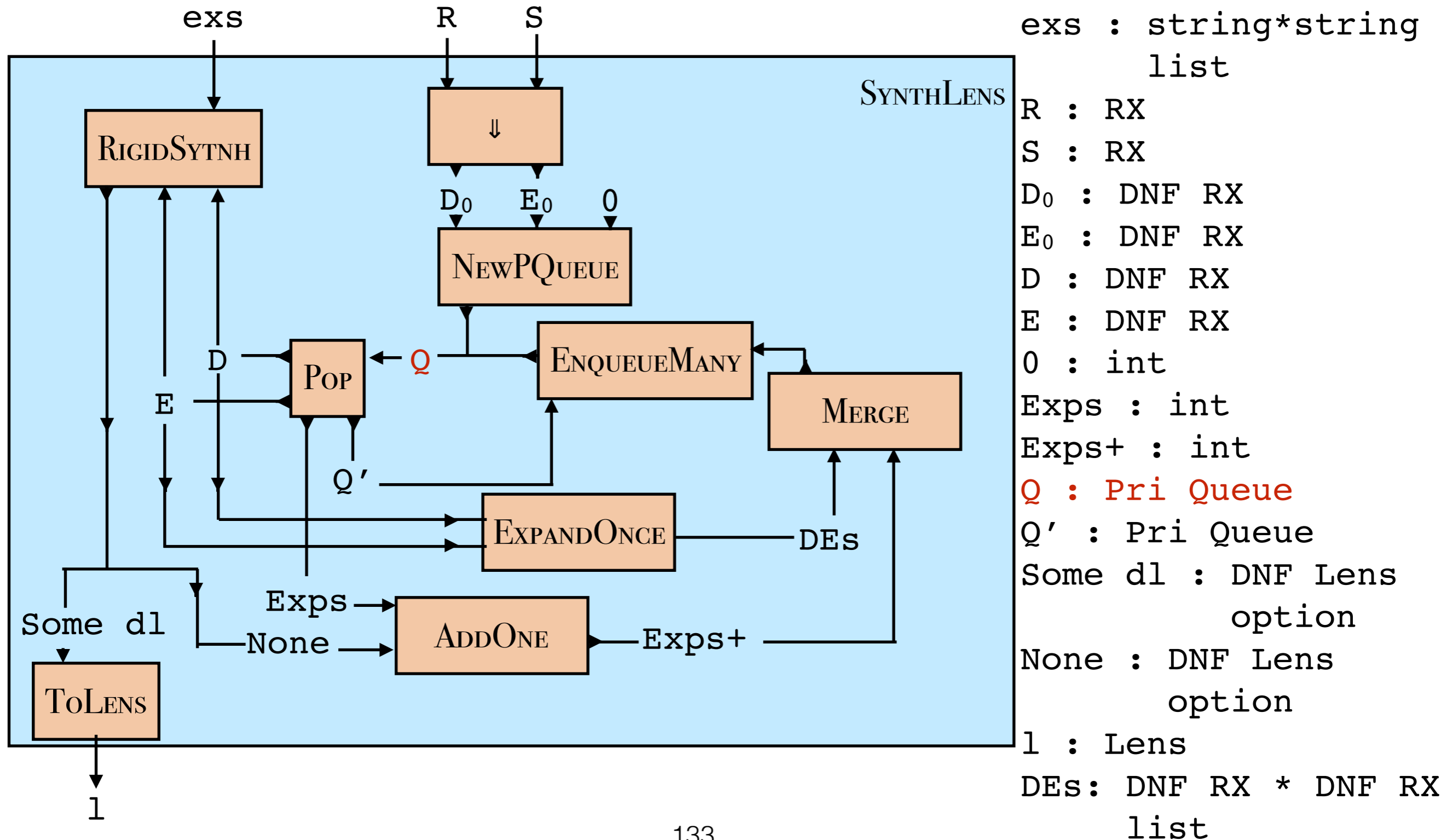
Synthesis Diagram



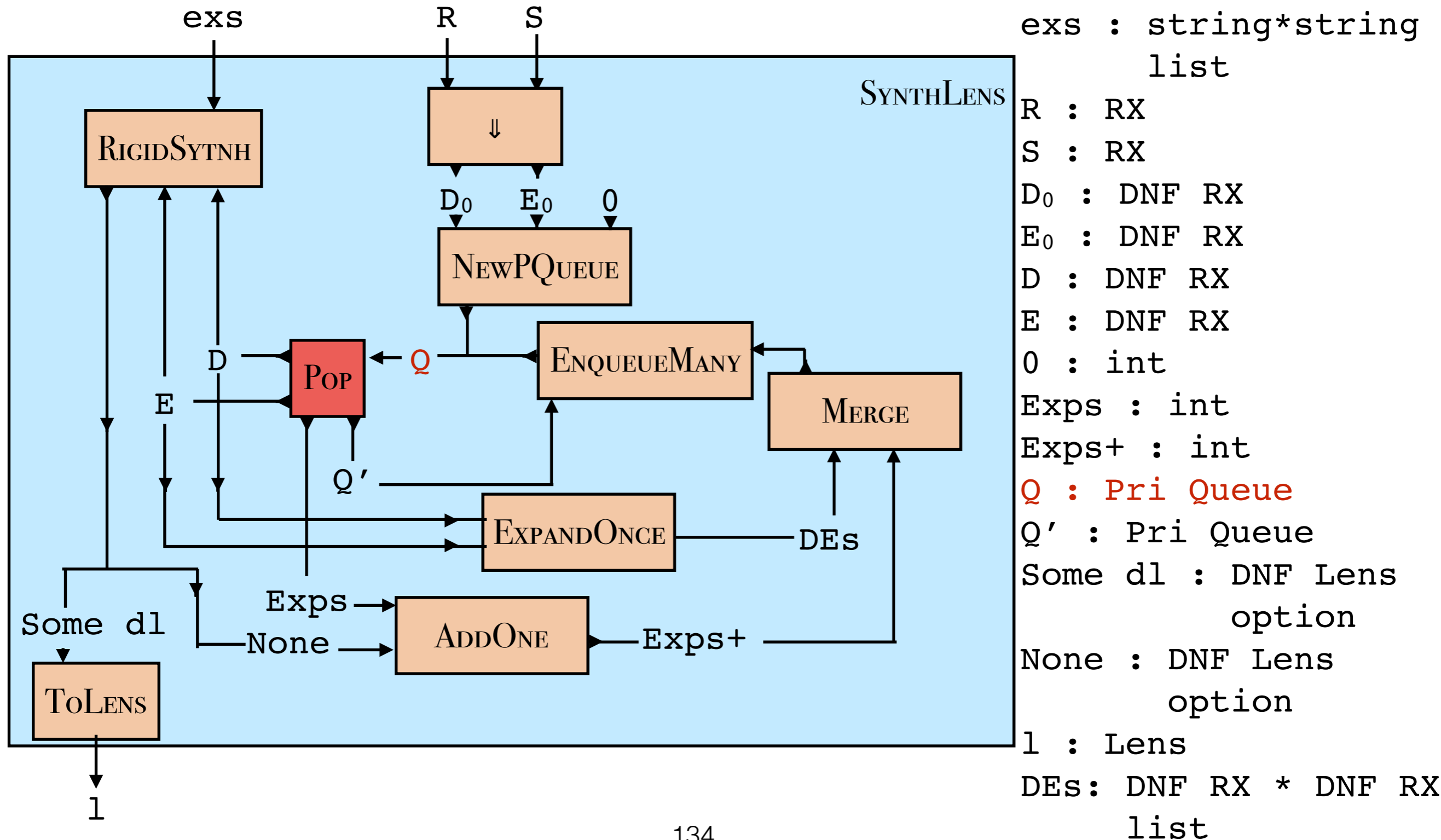
Synthesis Diagram



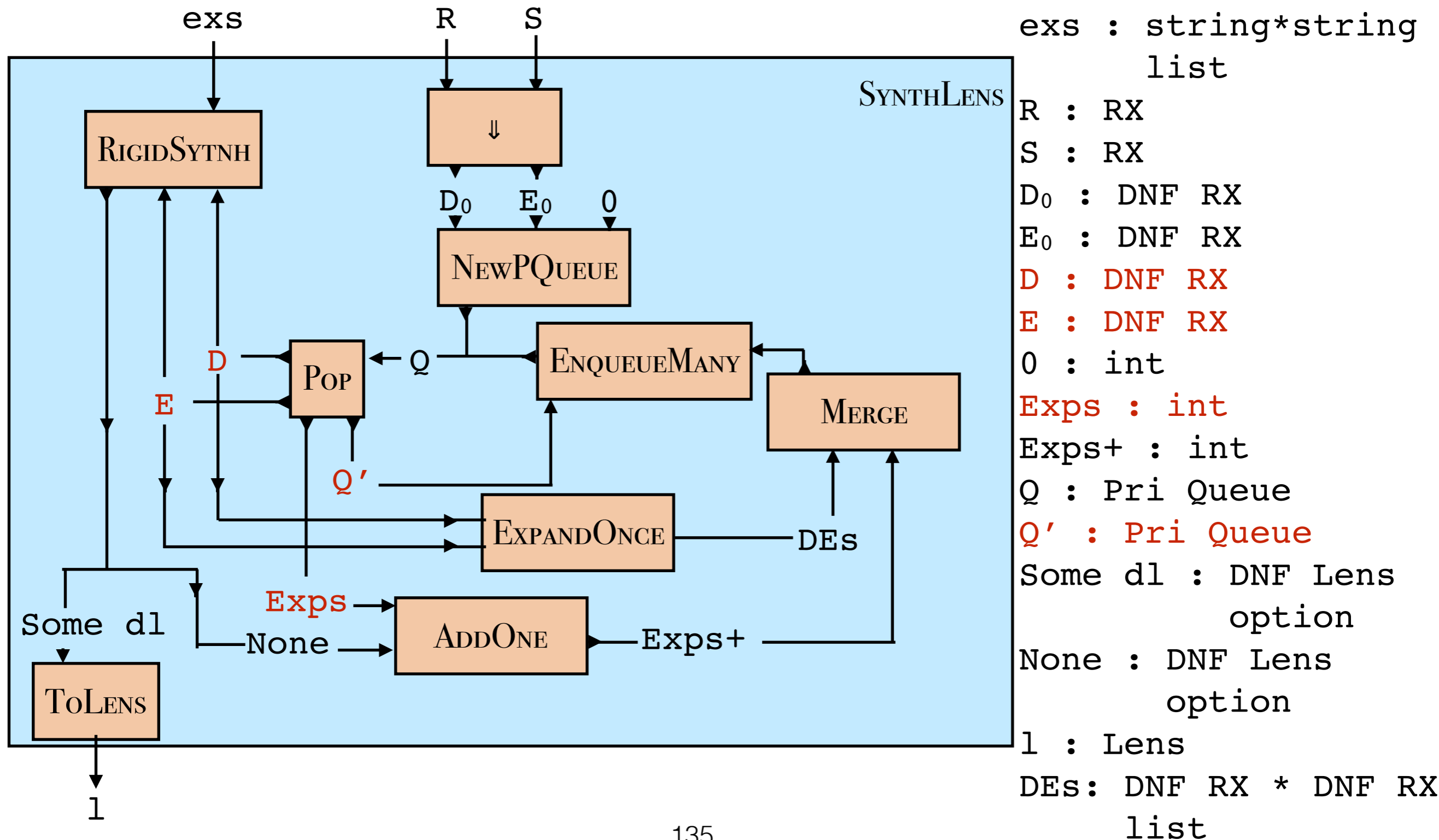
Synthesis Diagram



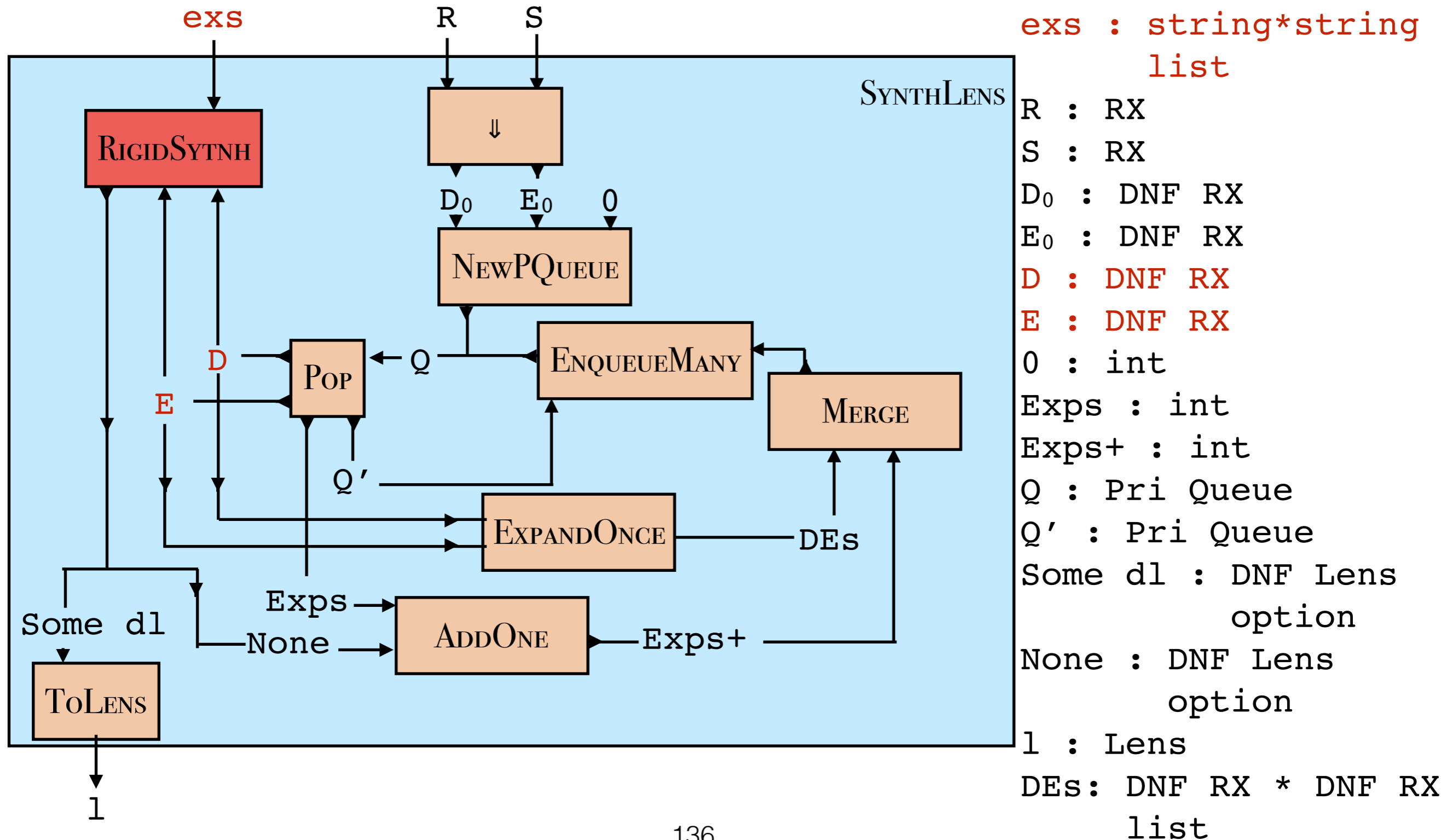
Synthesis Diagram



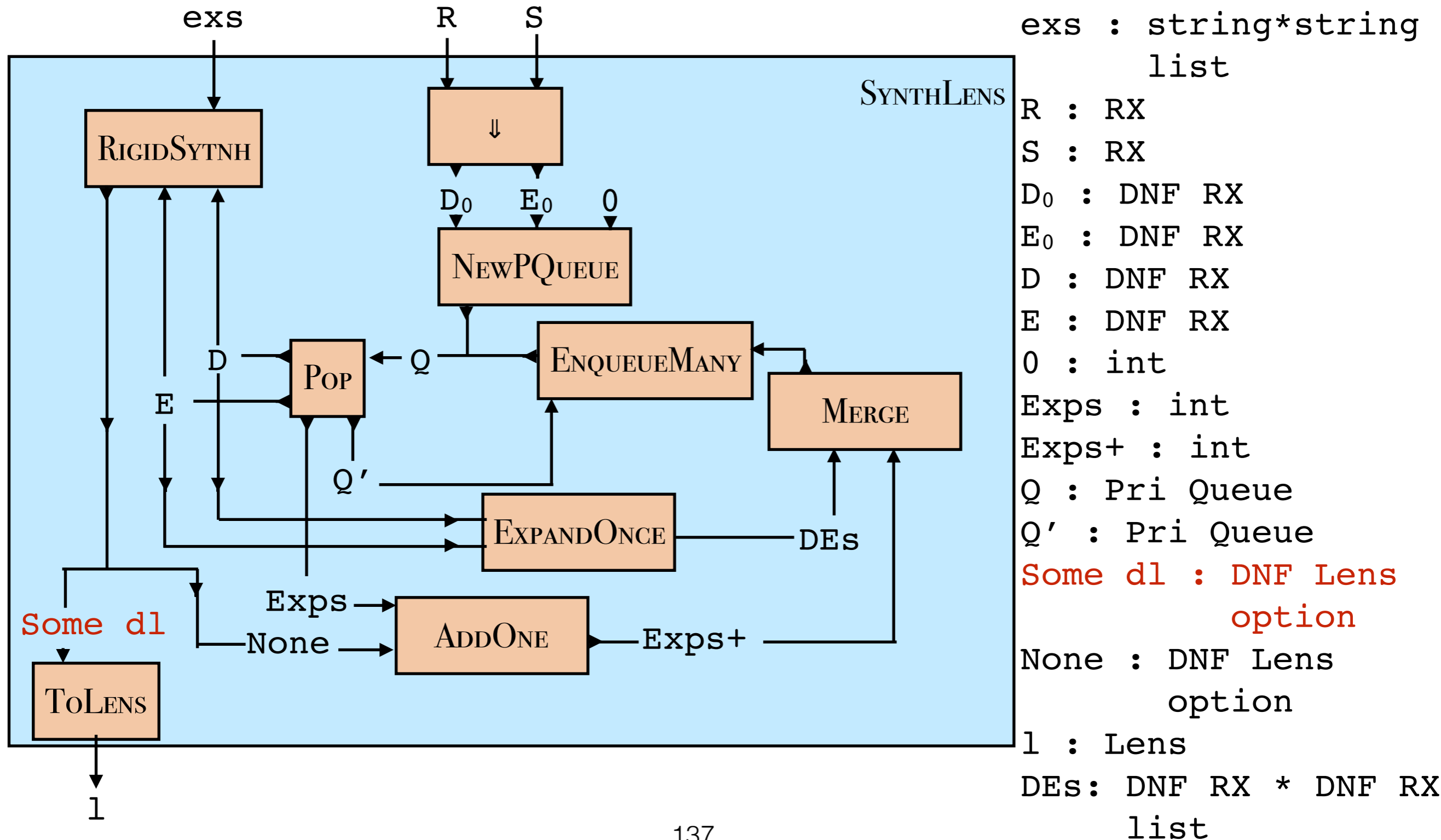
Synthesis Diagram



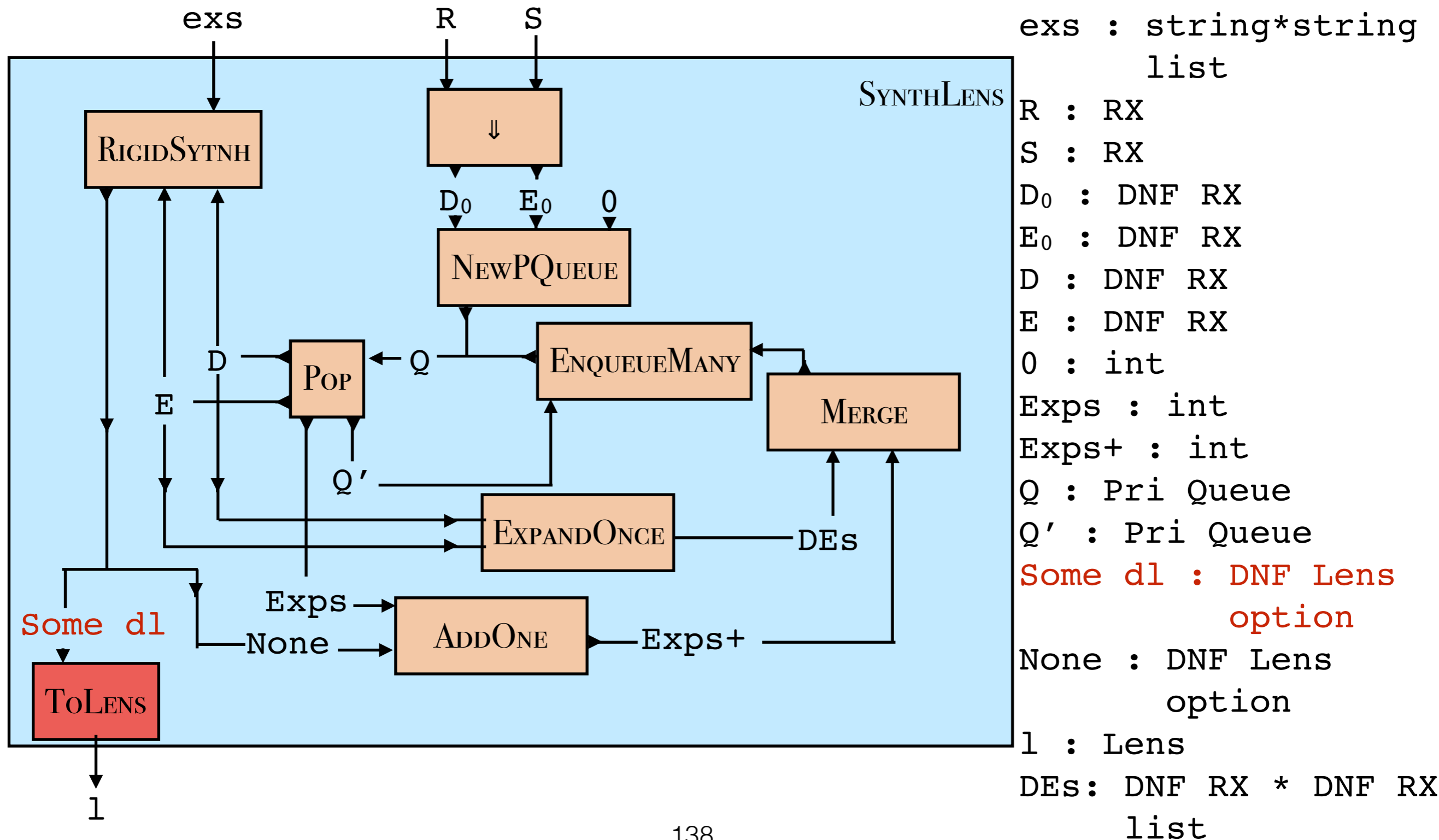
Synthesis Diagram



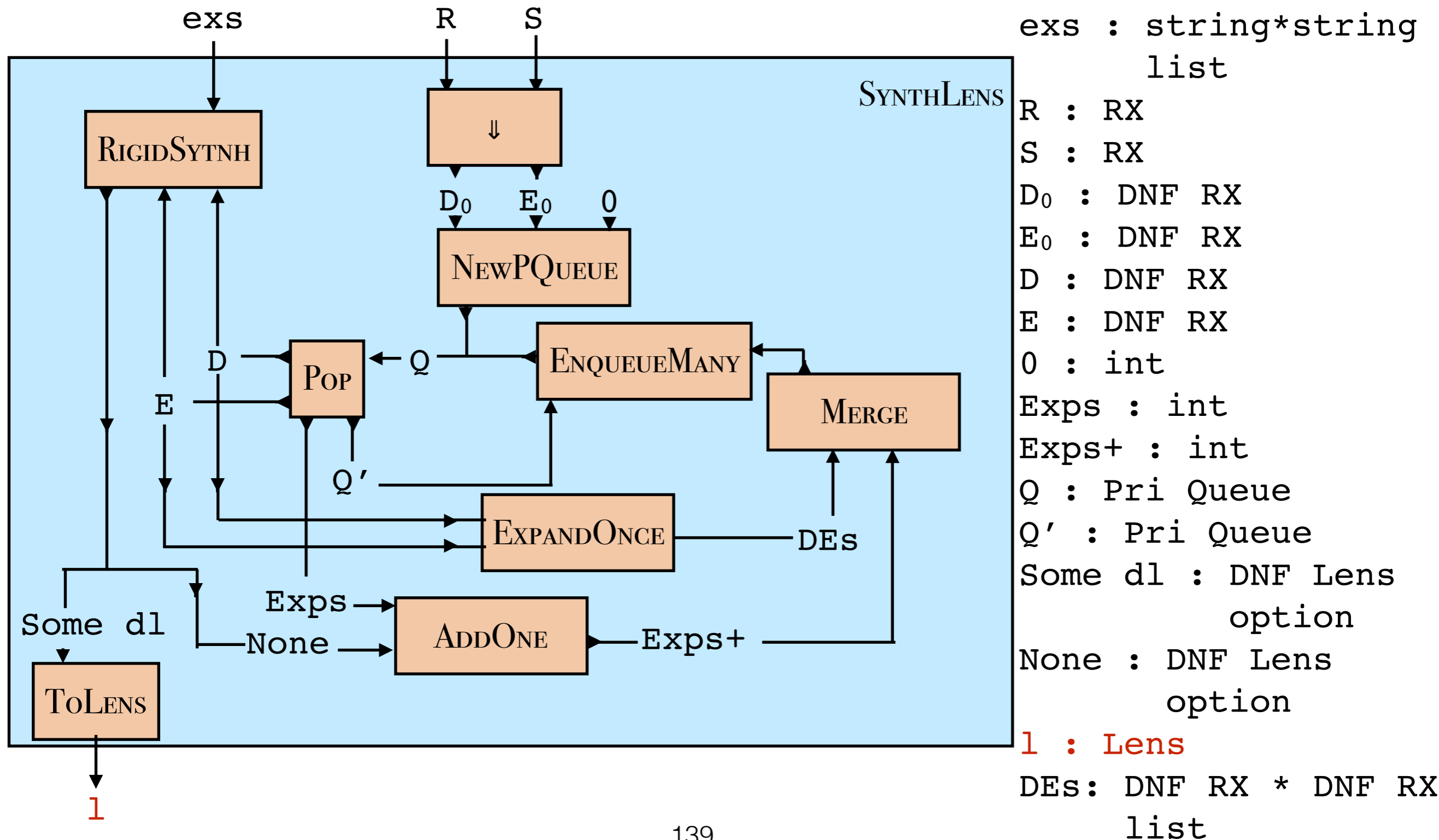
Synthesis Diagram



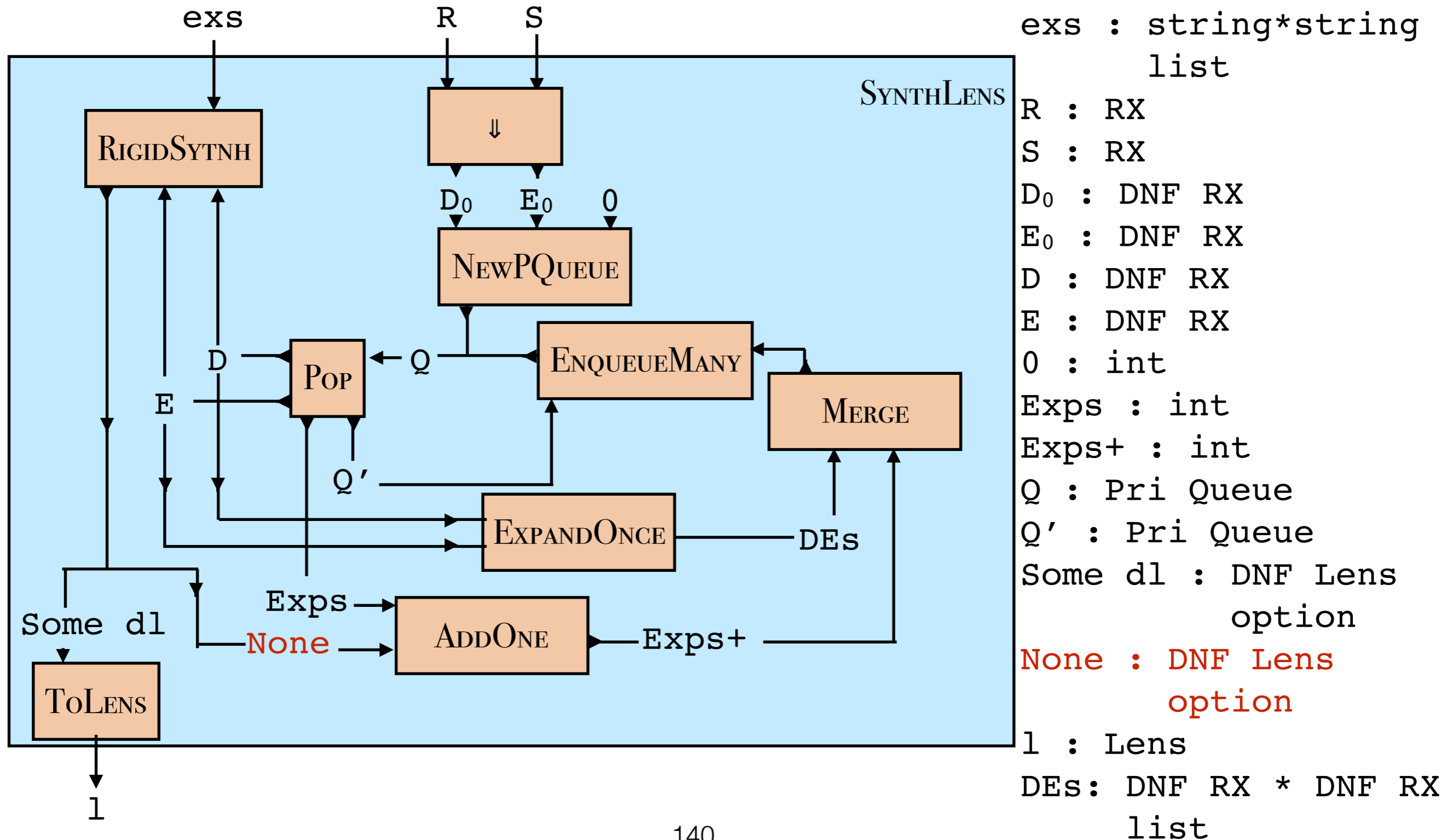
Synthesis Diagram



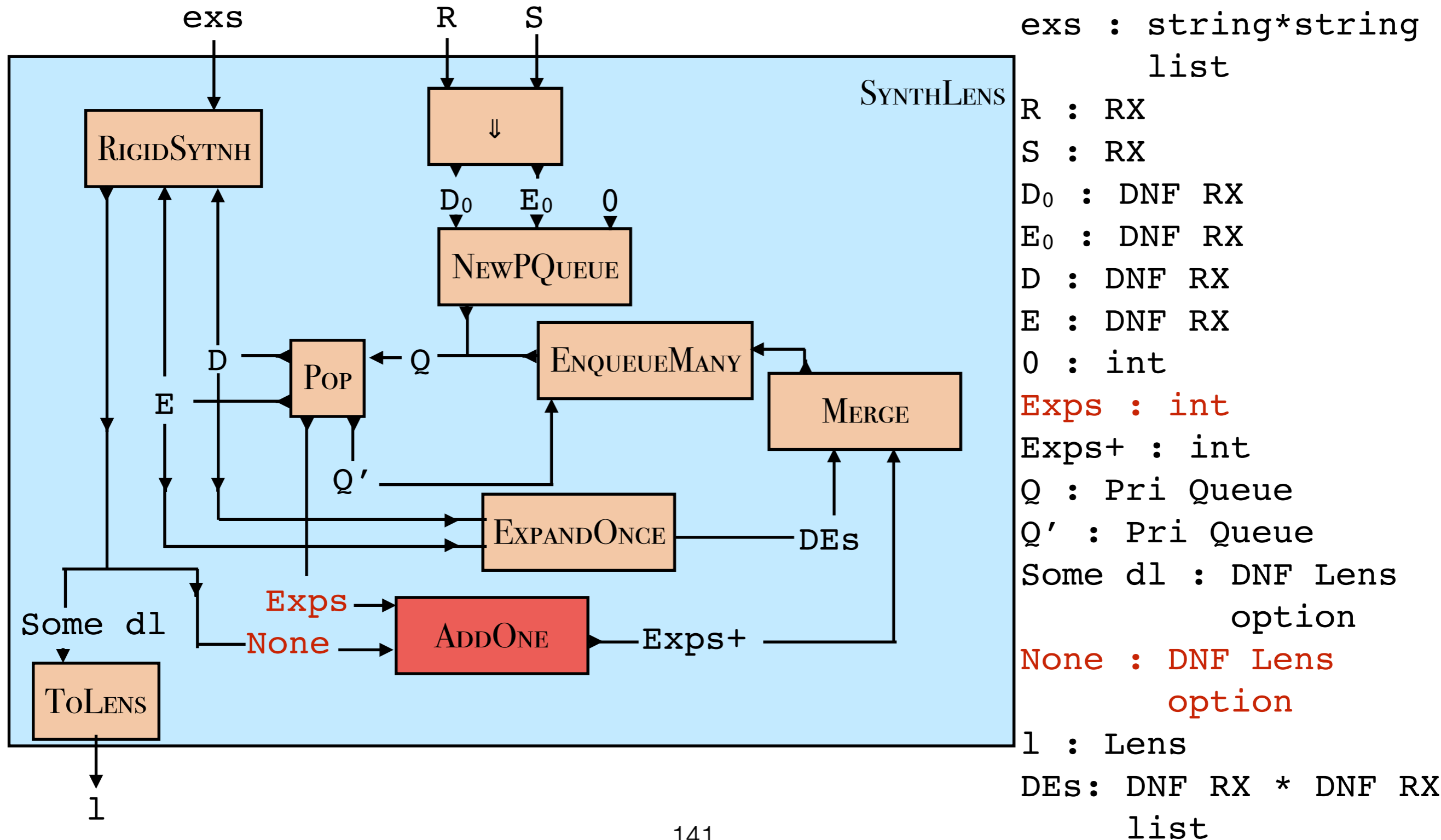
Synthesis Diagram



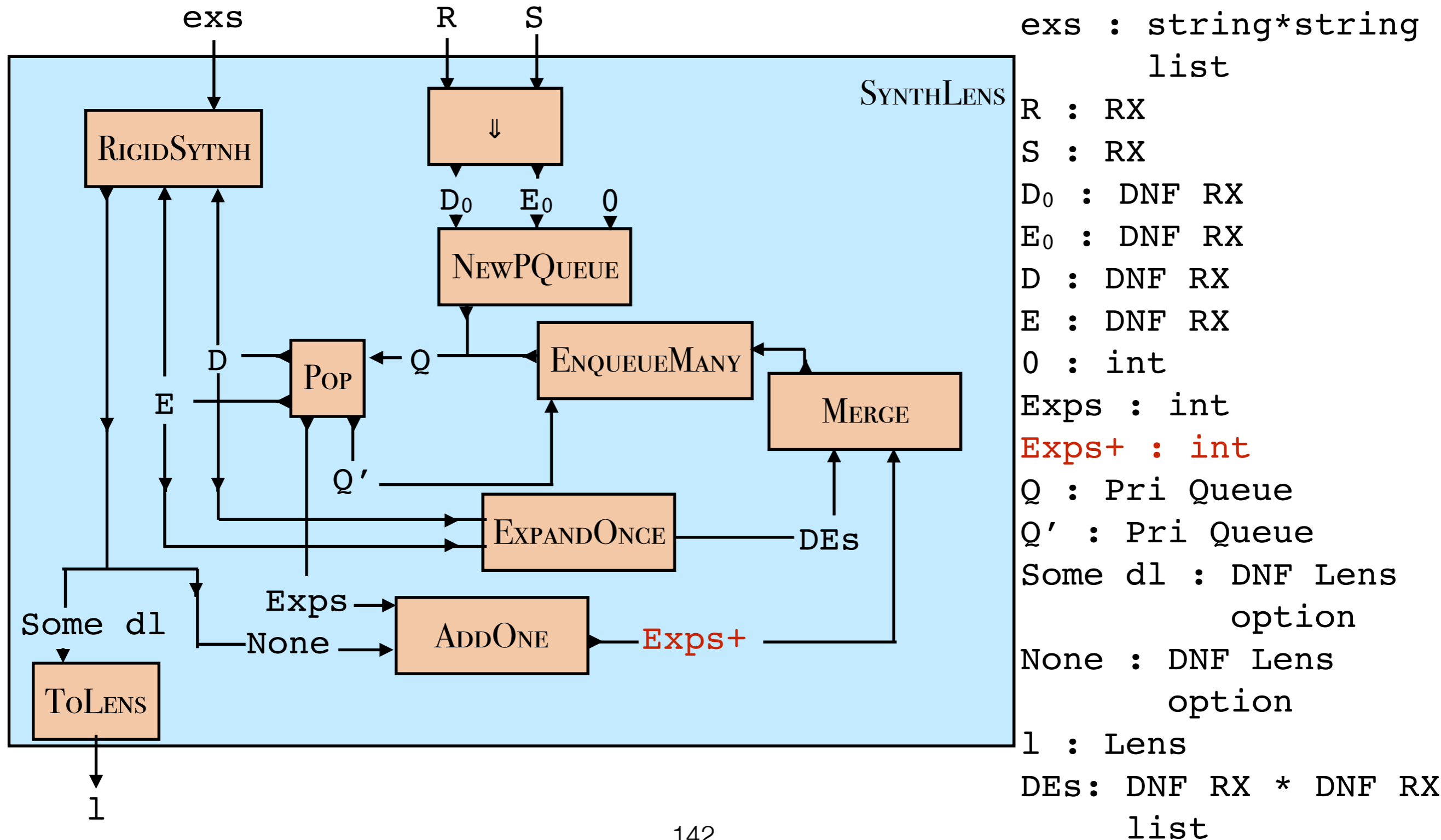
Synthesis Diagram



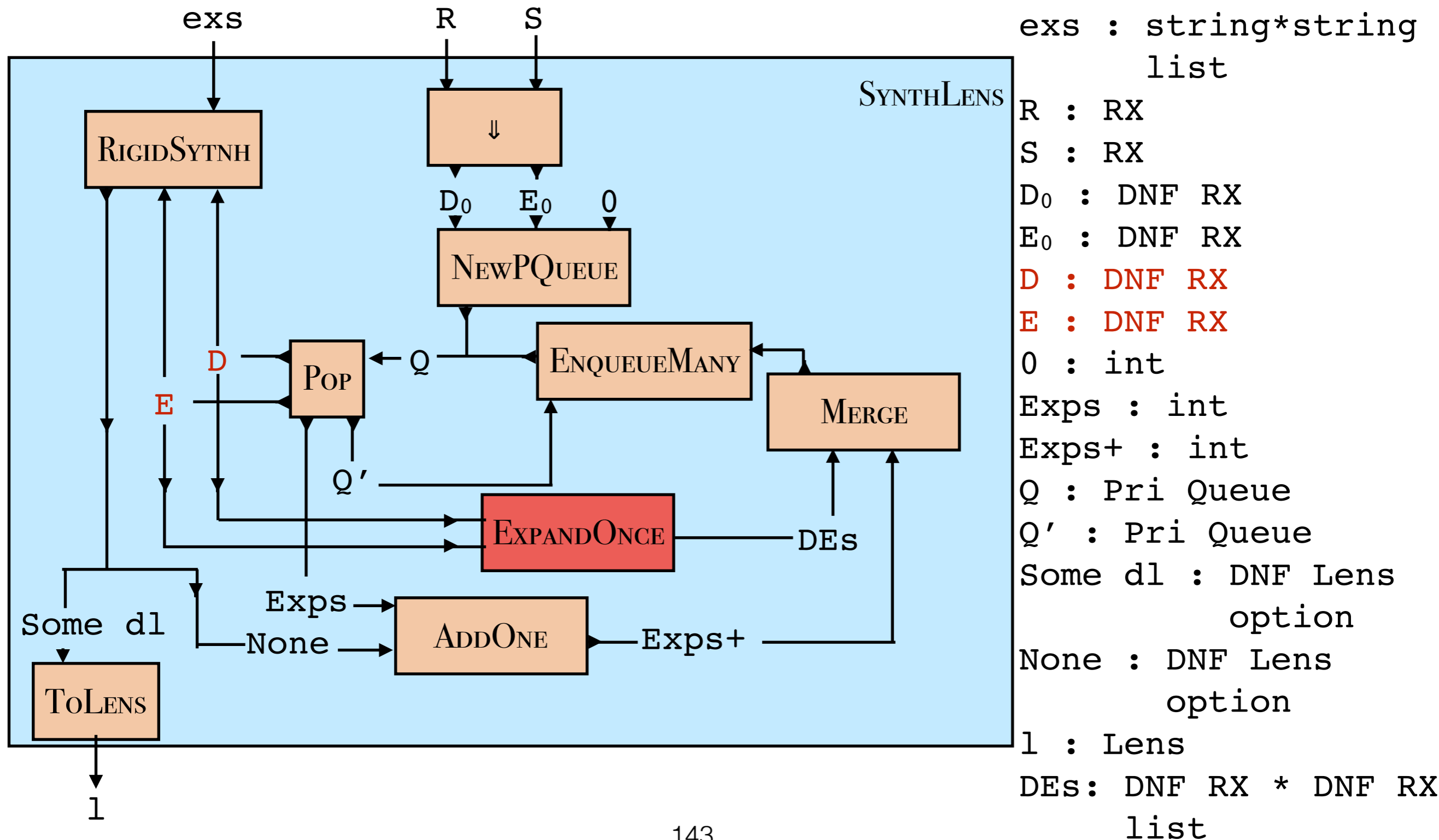
Synthesis Diagram



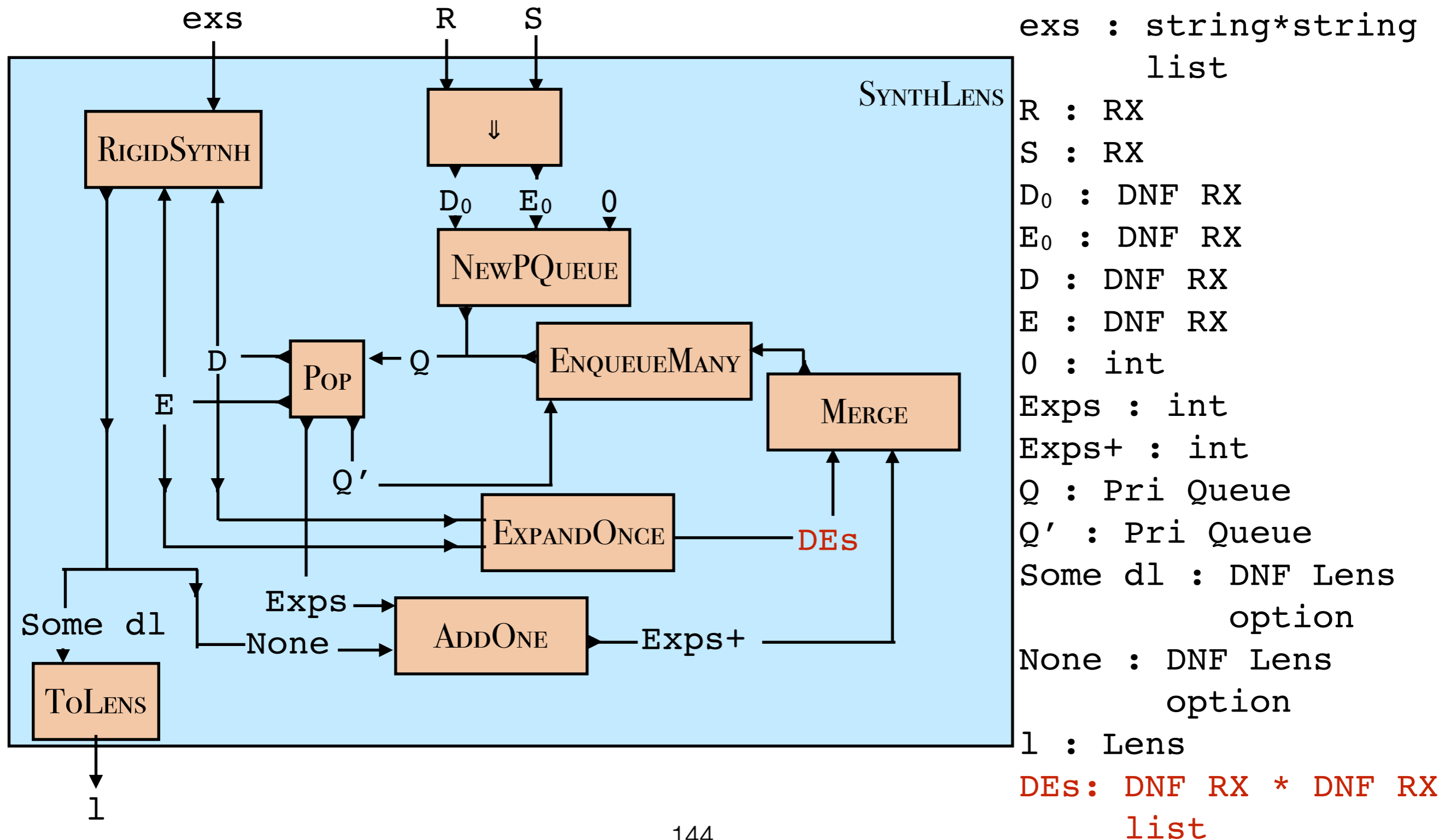
Synthesis Diagram



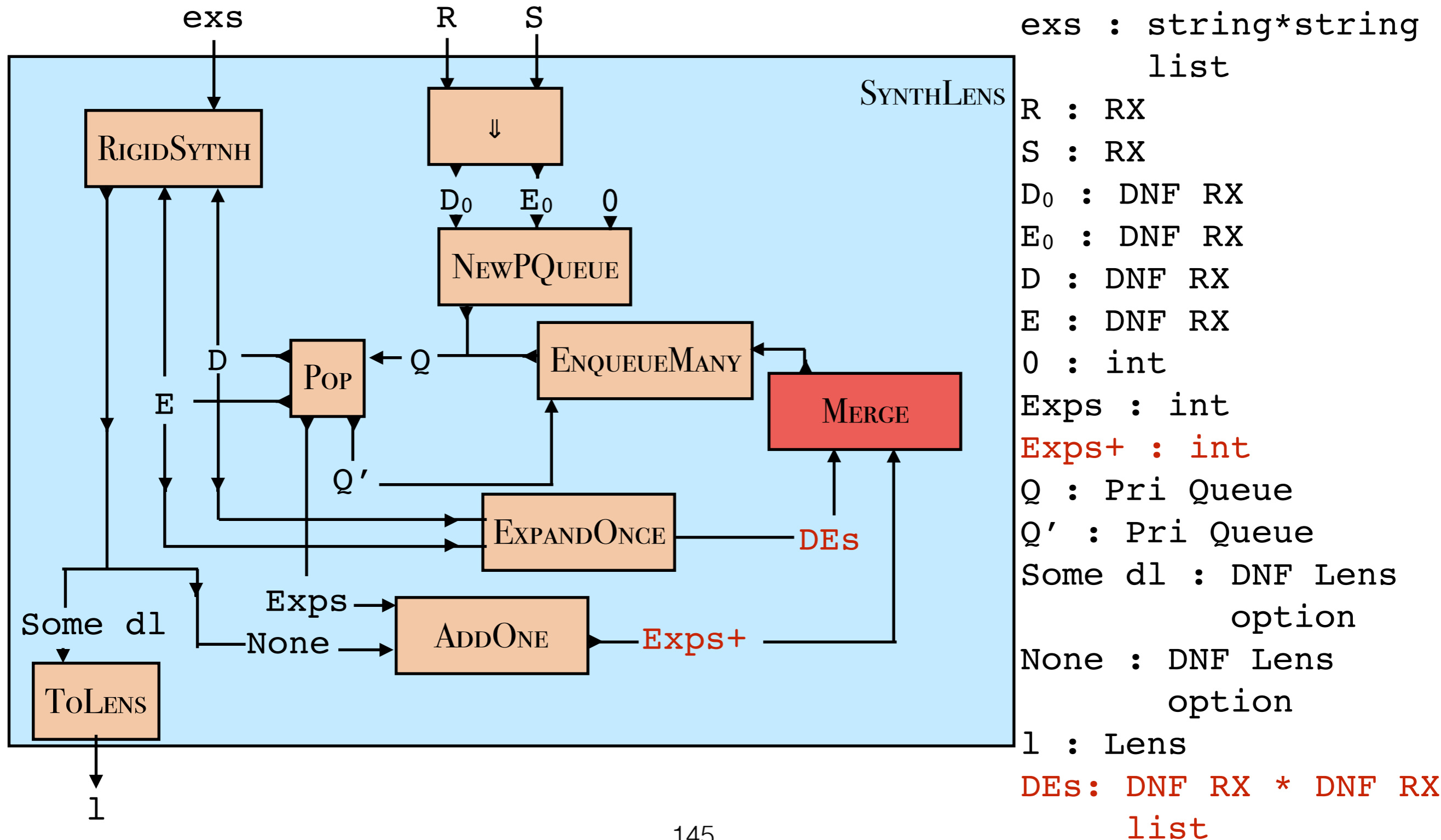
Synthesis Diagram



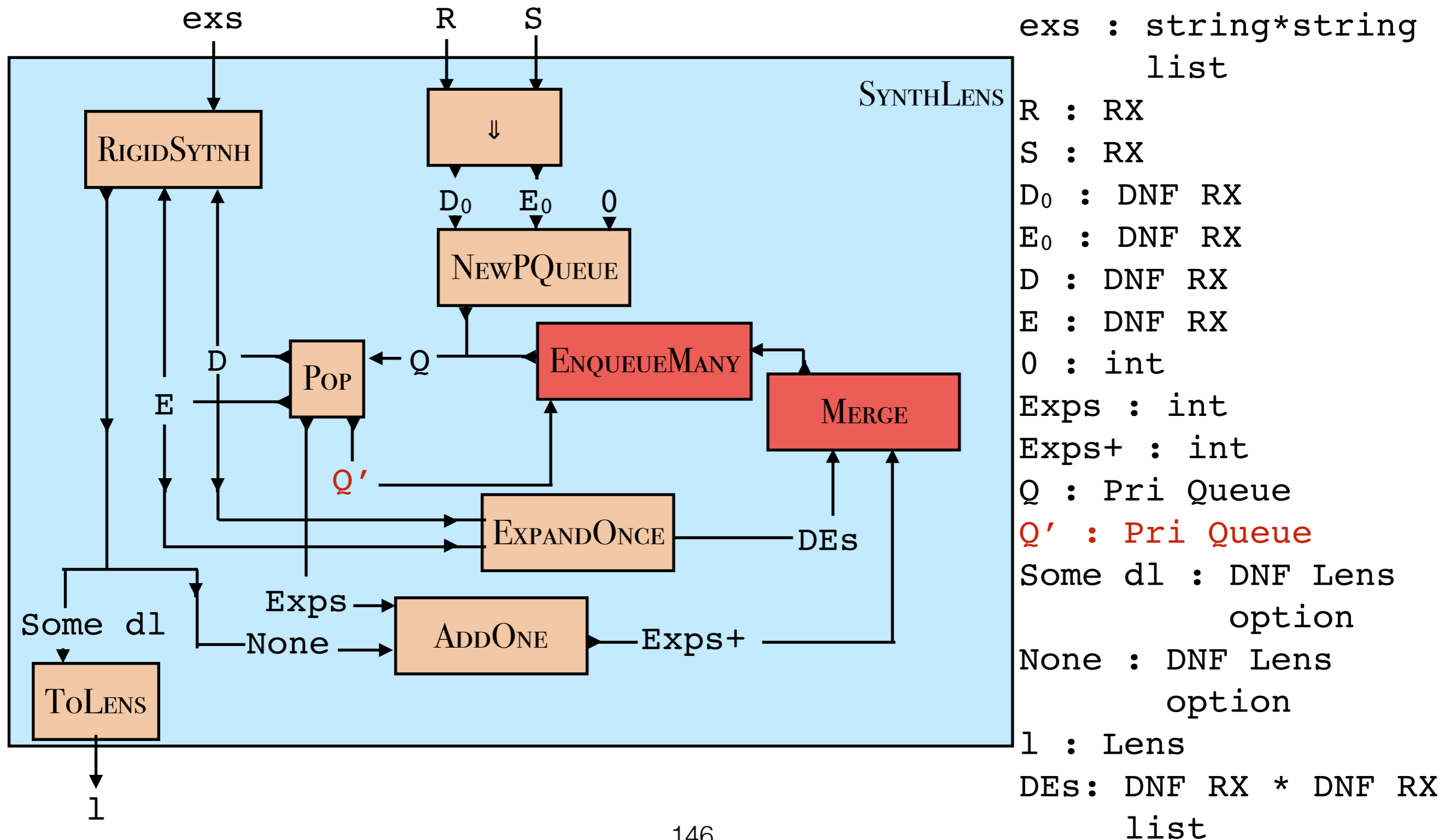
Synthesis Diagram



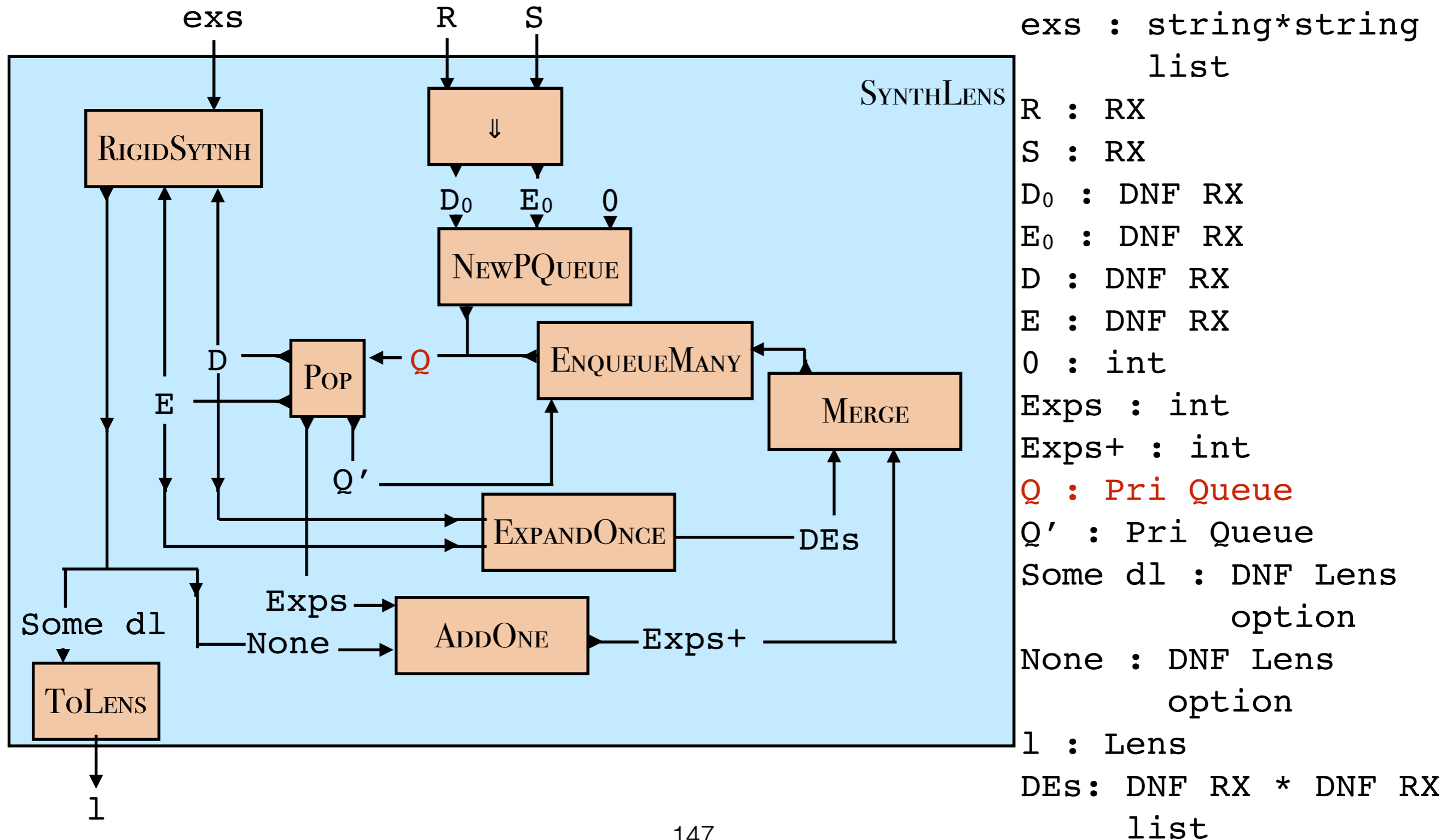
Synthesis Diagram



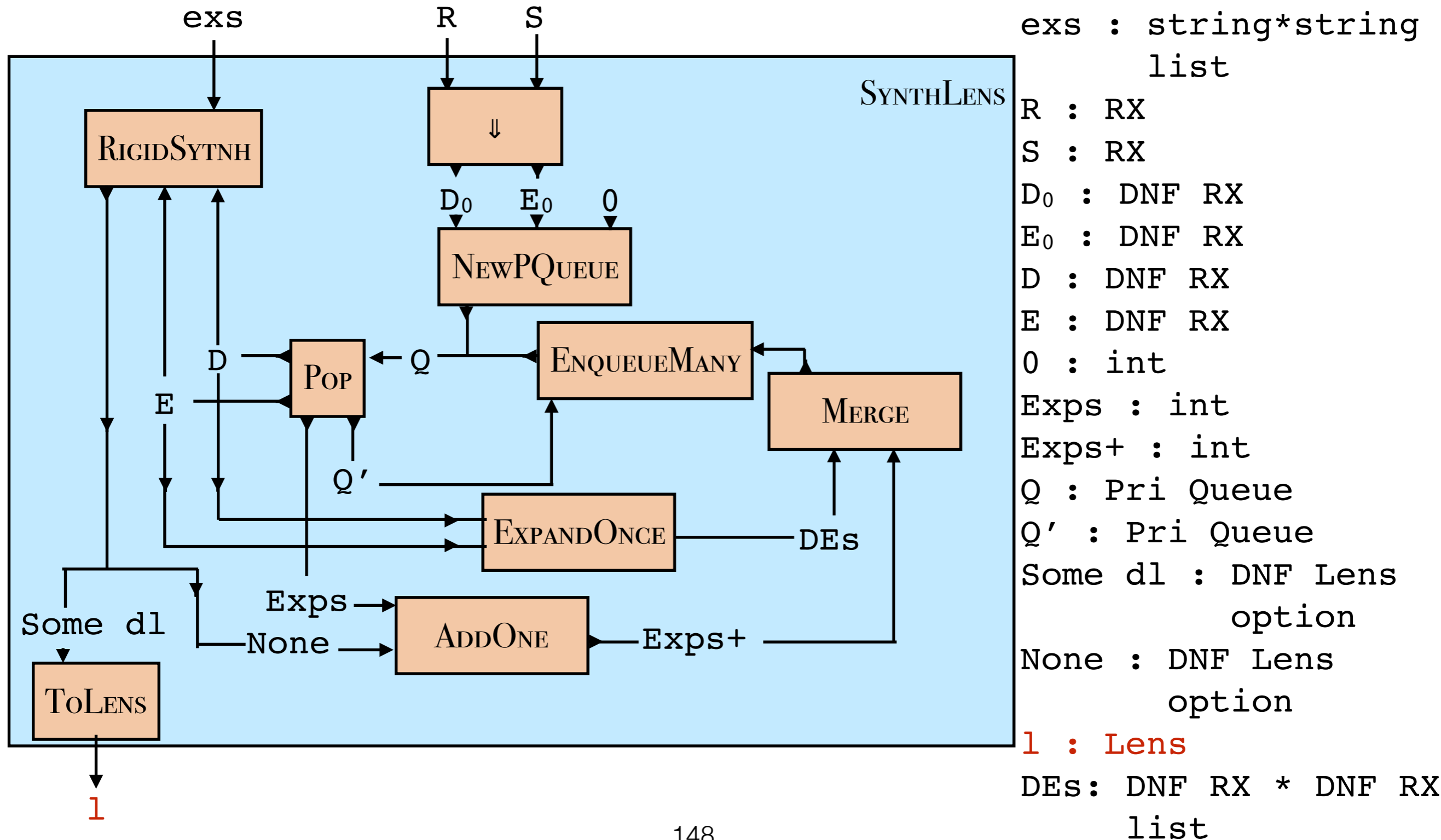
Synthesis Diagram



Synthesis Diagram



Synthesis Diagram



Priority Queue

- Higher priority for fewer expansions preformed
- Higher priority for more similar looking regular expressions
- Lower score means higher priority

$$\text{SCORE}(D,E,\text{Exps}) = d(D,E) + 8 * \text{Exps}$$

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry, 0) -> 1
```

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(Entry, 0) -> 1
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DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry, 0) -> 1  
(Names, 0) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry, 0) -> 1  
(Names, 0) -> 1  
(Names, 1) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry, 0) -> 1  
(Names, 0) -> 1  
(Names, 1) -> 1  
(Entry, 1) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry, 0) -> 1  
(Names, 0) -> 1  
(Names, 1) -> 1  
(Entry, 1) -> 1  
(Other, 0) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry, 0) -> 1  
(Names, 0) -> 1  
(Names, 1) -> 1  
(Entry, 1) -> 1  
(Other, 0) -> 1
```

```
(Names, 0) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry,0) -> 1  
(Names,0) -> 1  
(Names,1) -> 1  
(Entry,1) -> 1  
(Other,0) -> 1
```

```
(Names,0) -> 1  
(Alt,1) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry,0) -> 1  
(Names,0) -> 1  
(Names,1) -> 1  
(Entry,1) -> 1  
(Other,0) -> 1
```

```
(Names,0) -> 1  
(Alt,1) -> 1  
(Entry,0) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

```
(Entry,0) -> 1  
(Names,0) -> 1  
(Names,1) -> 1  
(Entry,1) -> 1  
(Other,0) -> 1
```

```
(Names,0) -> 2  
(Alt,1) -> 1  
(Entry,0) -> 1
```


DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

```
Entry."test" | Names | (Names.Entry)* | Other  
Names | Alt*.Entry | Names.Names
```

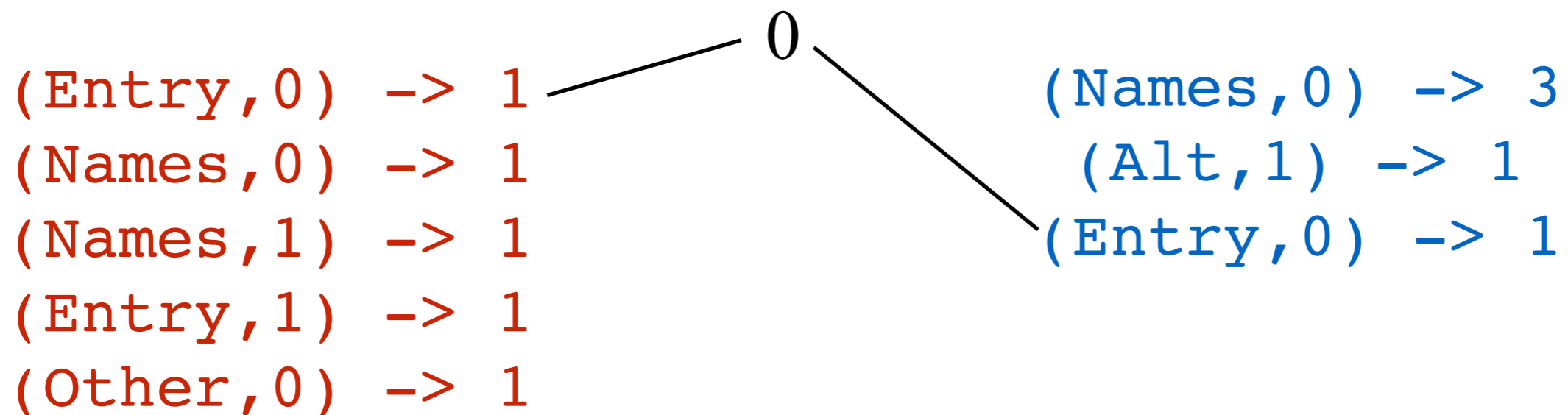
```
(Entry,0) -> 1  
(Names,0) -> 1  
(Names,1) -> 1  
(Entry,1) -> 1  
(Other,0) -> 1
```

```
(Names,0) -> 3  
(Alt,1) -> 1  
(Entry,0) -> 1
```

DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

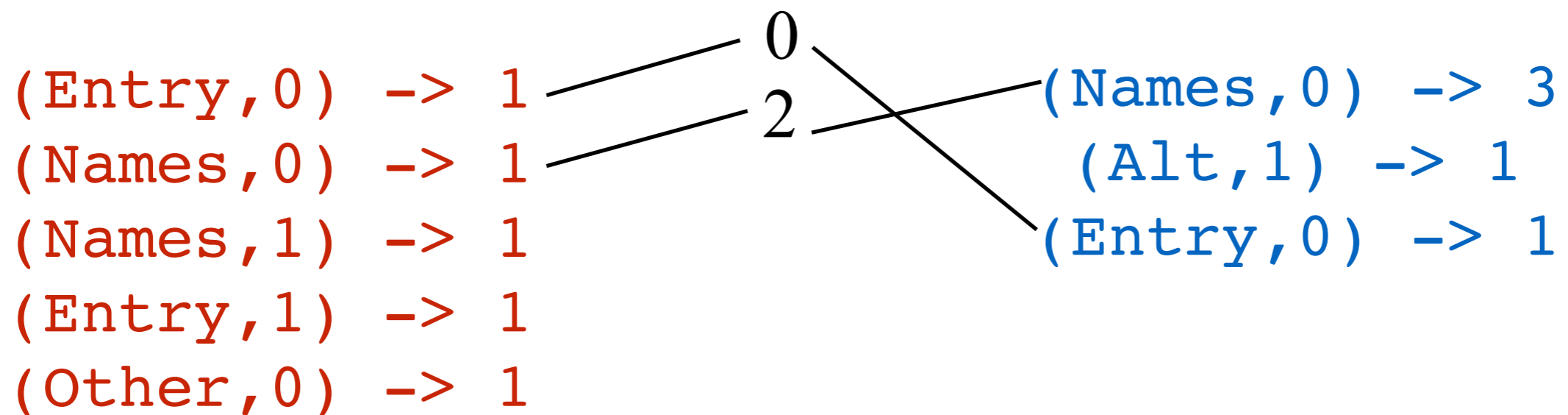
Entry."test" | Names | (Names.Entry)* | Other
Names | Alt*.Entry | Names.Names



DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

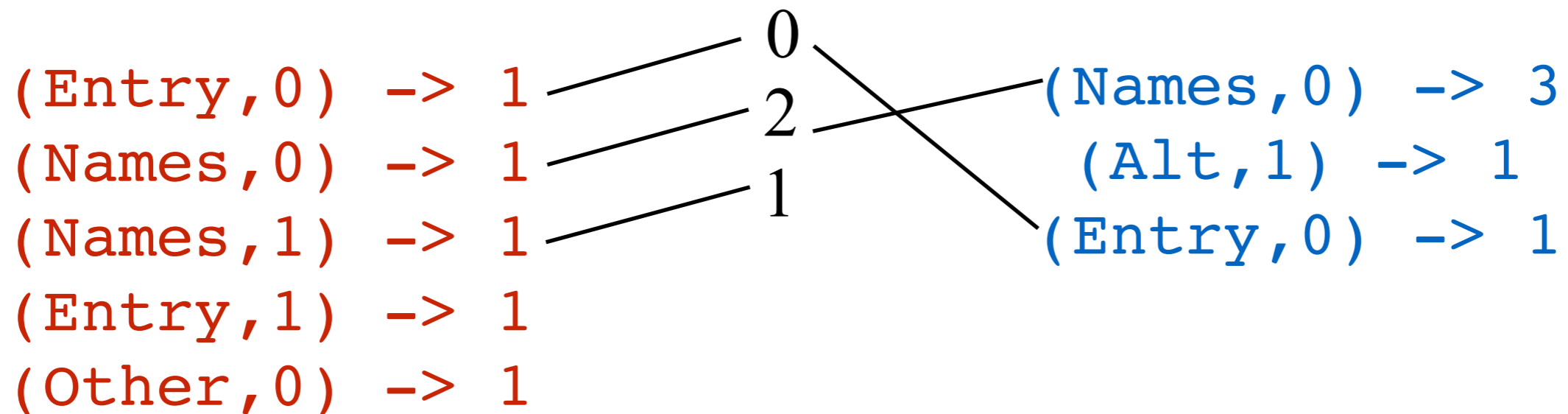
Entry."test" | Names | (Names.Entry)* | Other
Names | Alt*.Entry | Names.Names



DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

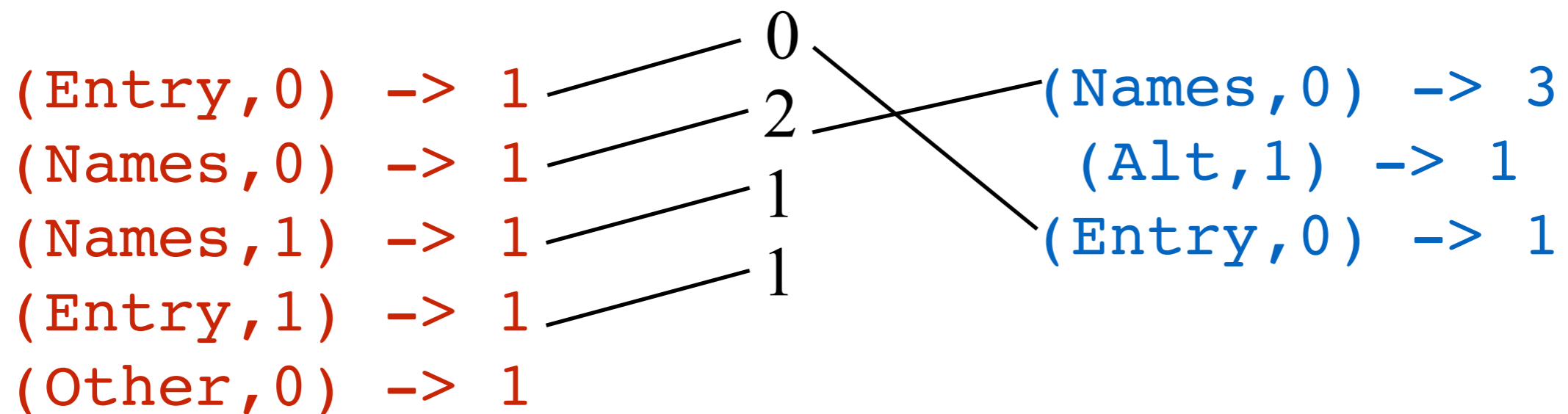
Entry."test" | Names | (Names.Entry)* | Other
Names | Alt*.Entry | Names.Names



DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

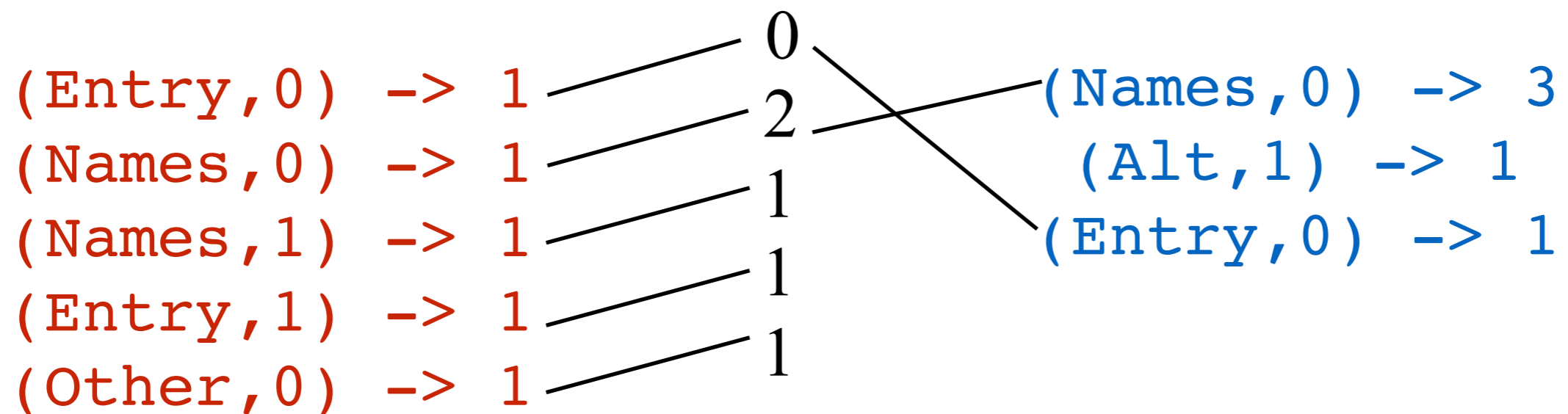
Entry."test" | Names | (Names.Entry)* | Other
Names | Alt*.Entry | Names.Names



DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

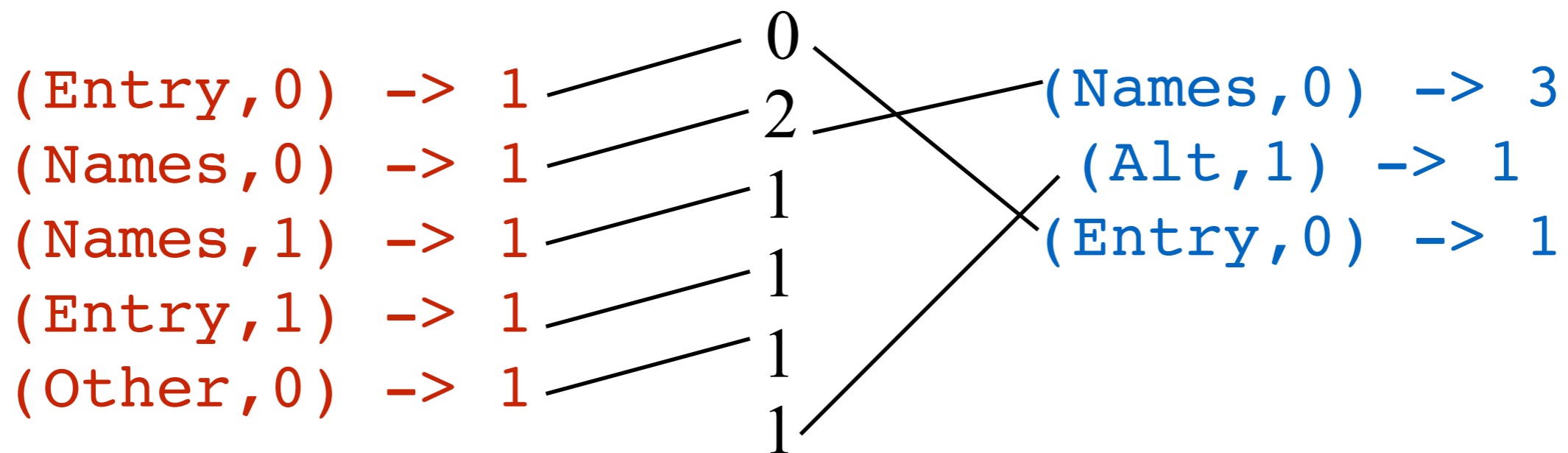
Entry."test" | Names | (Names.Entry)* | Other
Names | Alt*.Entry | Names.Names



DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

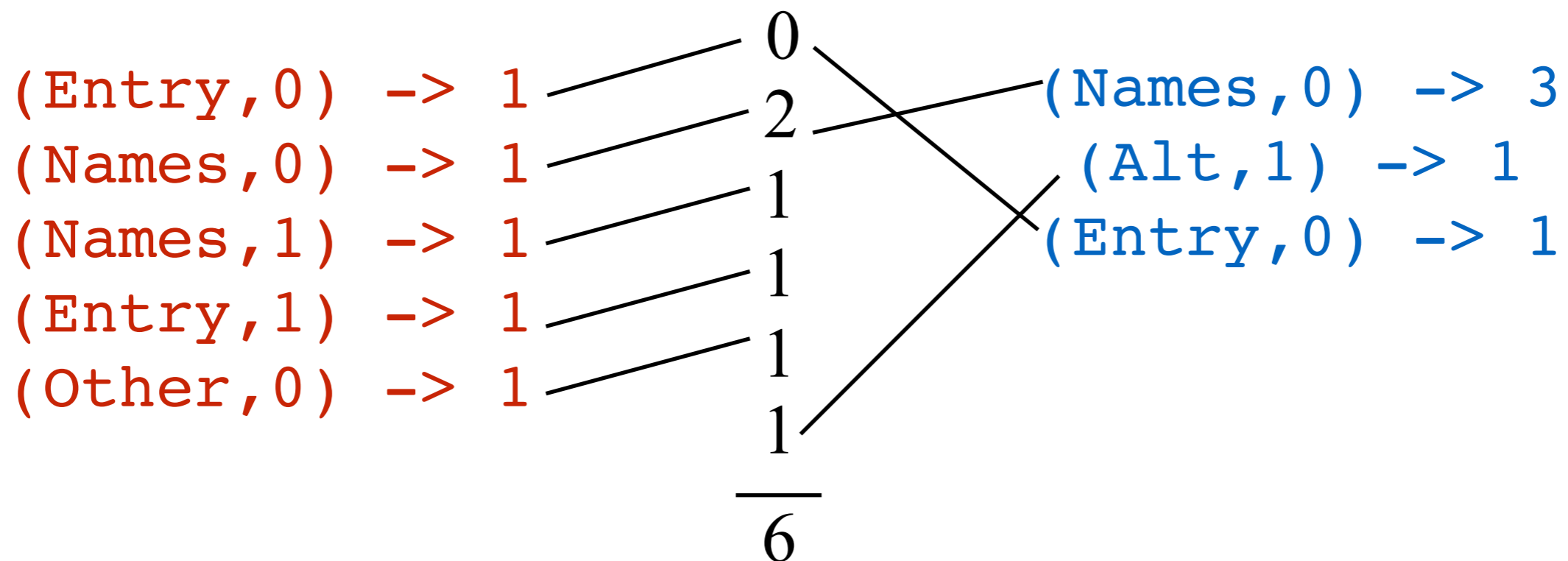
Entry."test" | Names | (Names.Entry)* | Other
Names | Alt*.Entry | Names.Names



DNF RX Distance

Based on the distribution of the user-defined Regular Expressions

Entry."test" | Names | (Names.Entry)* | Other
Names | Alt*.Entry | Names.Names

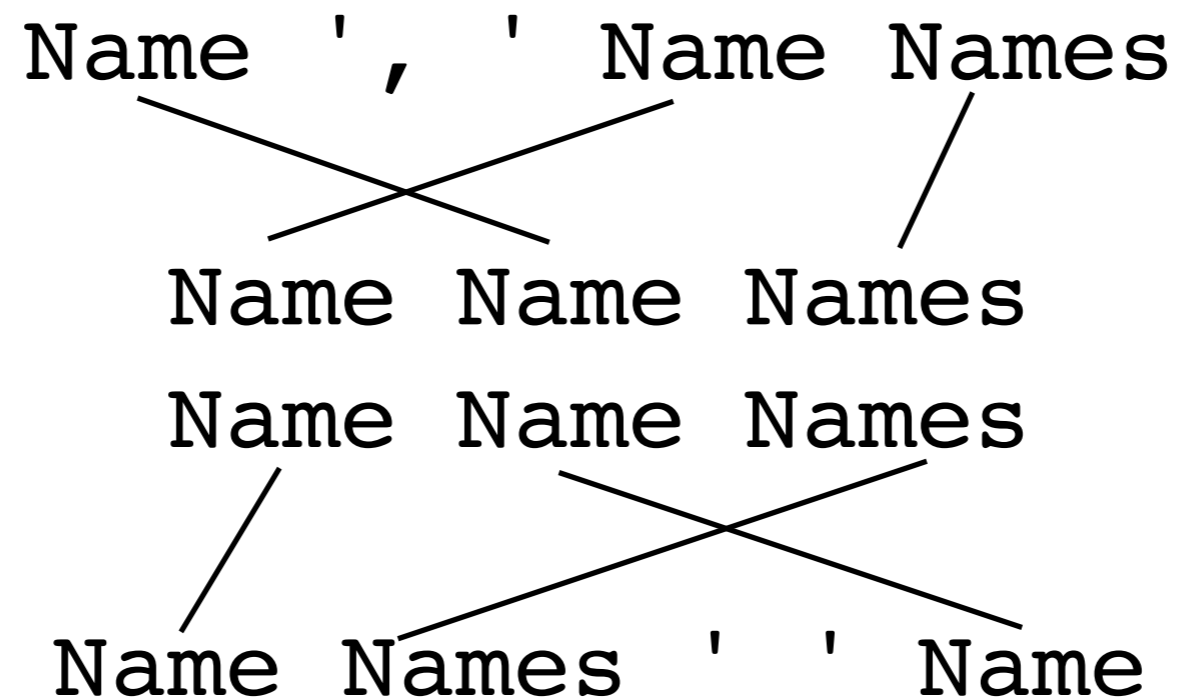


RIGIDSYNTH

Synthesize $d1$ satisfying $d1 :: DR \Leftrightarrow DS$

Difficulty lies in determining the permutations σ

Order the subcomponents of a DNF Regular Expression

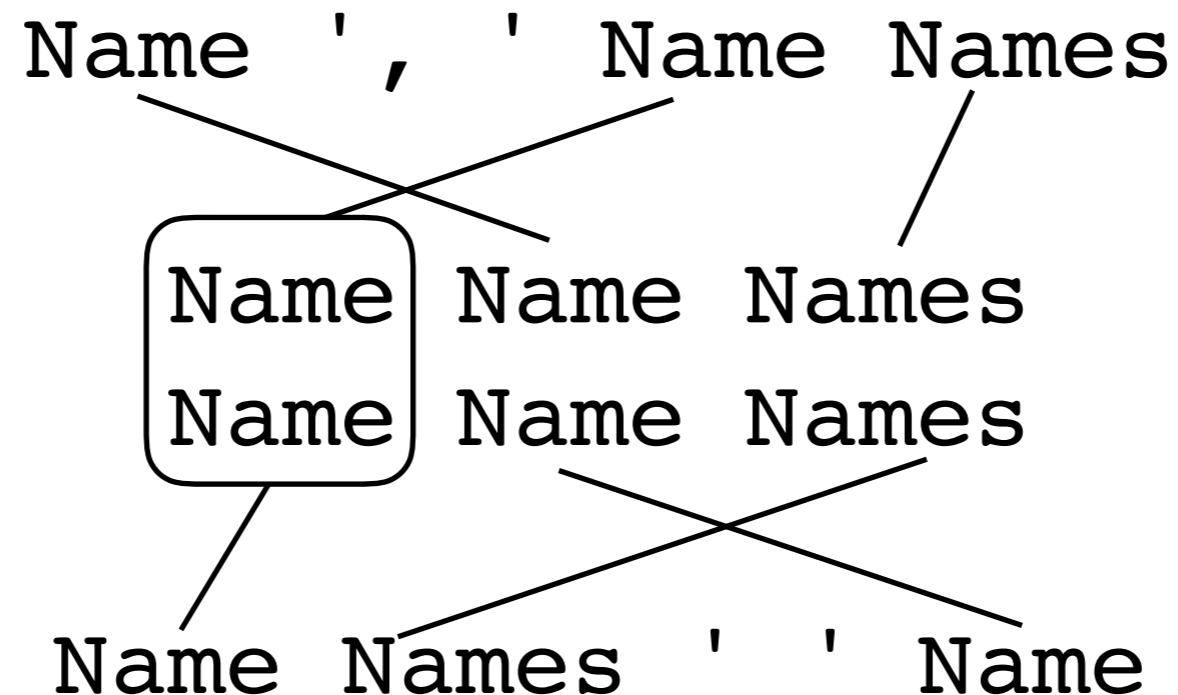


RIGIDSYNTH

Synthesize $d1$ satisfying $d1 :: DR \Leftrightarrow DS$

Difficulty lies in determining the permutations σ

Order the subcomponents of a DNF Regular Expression

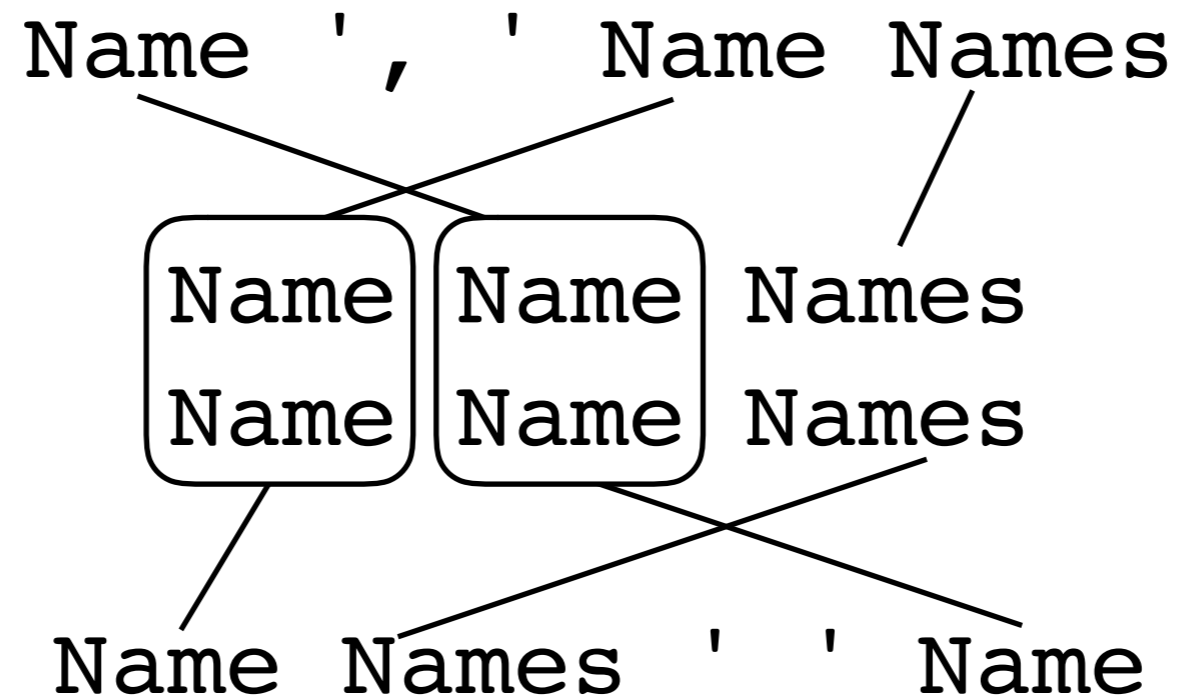


RIGIDSYNTH

Synthesize $d1$ satisfying $d1 :: DR \Leftrightarrow DS$

Difficulty lies in determining the permutations σ

Order the subcomponents of a DNF Regular Expression

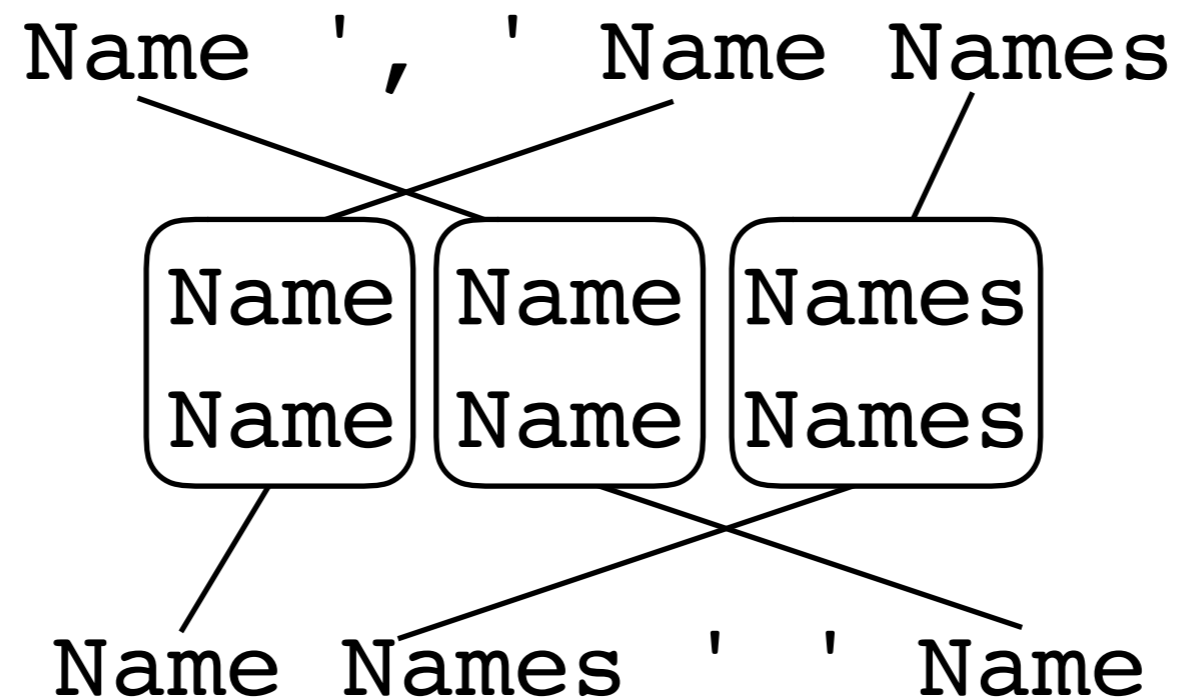


RIGIDSYNTH

Synthesize $d1$ satisfying $d1 :: DR \Leftrightarrow DS$

Difficulty lies in determining the permutations σ

Order the subcomponents of a DNF Regular Expression

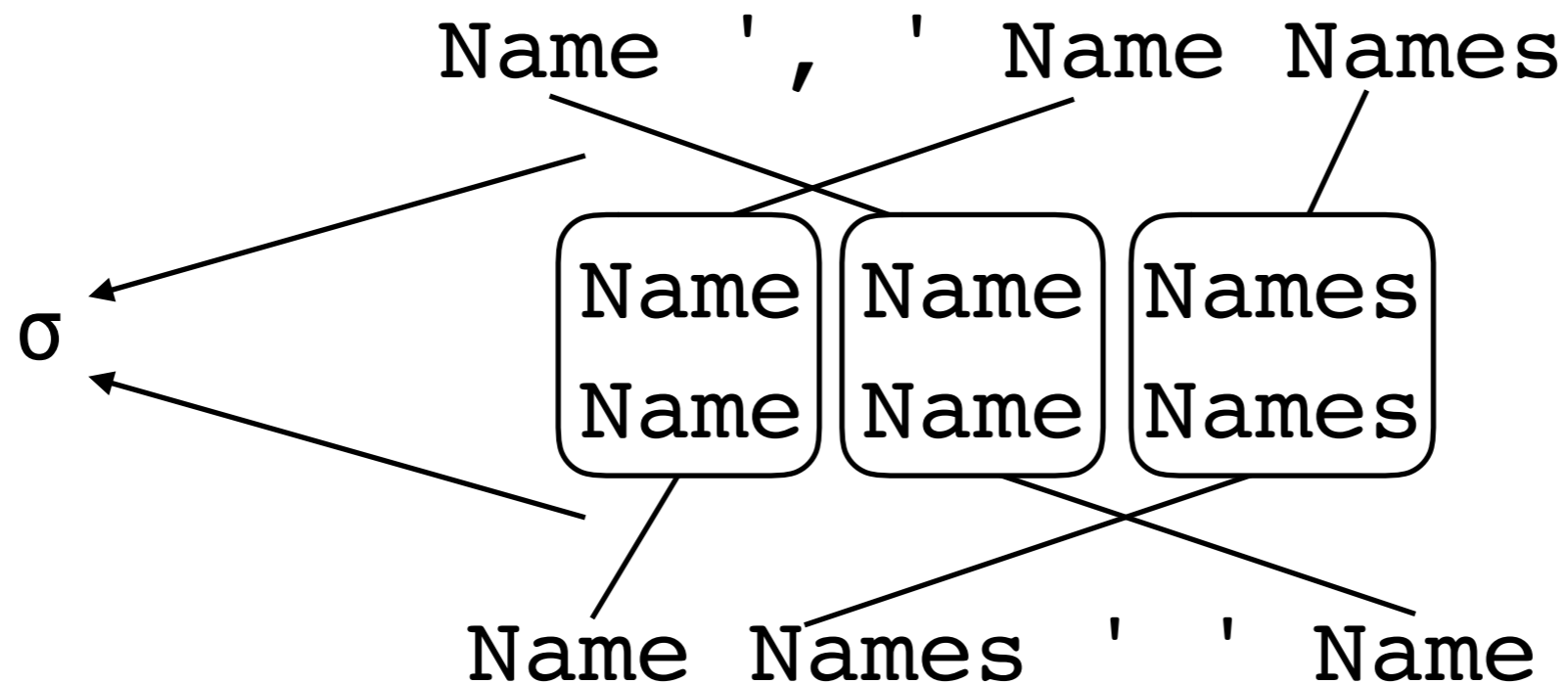


RIGIDSYNTH

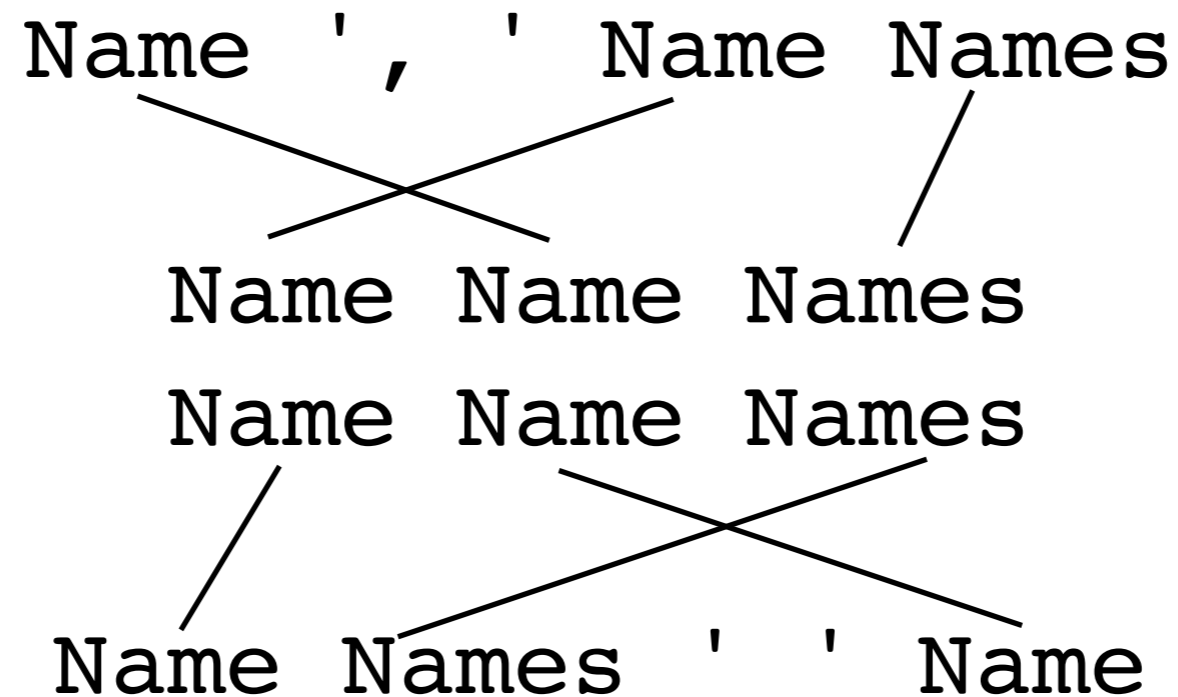
Synthesize $d1$ satisfying $d1 :: DR \Leftrightarrow DS$

Difficulty lies in determining the permutations σ

Order the subcomponents of a DNF Regular Expression



Satisfying Examples



Satisfying Examples

Add in additional information from the parsing of input/
output example strings!

```
Name [ Jones ] ' , ' Name [ Simon ] Names [ Peyton ]
```

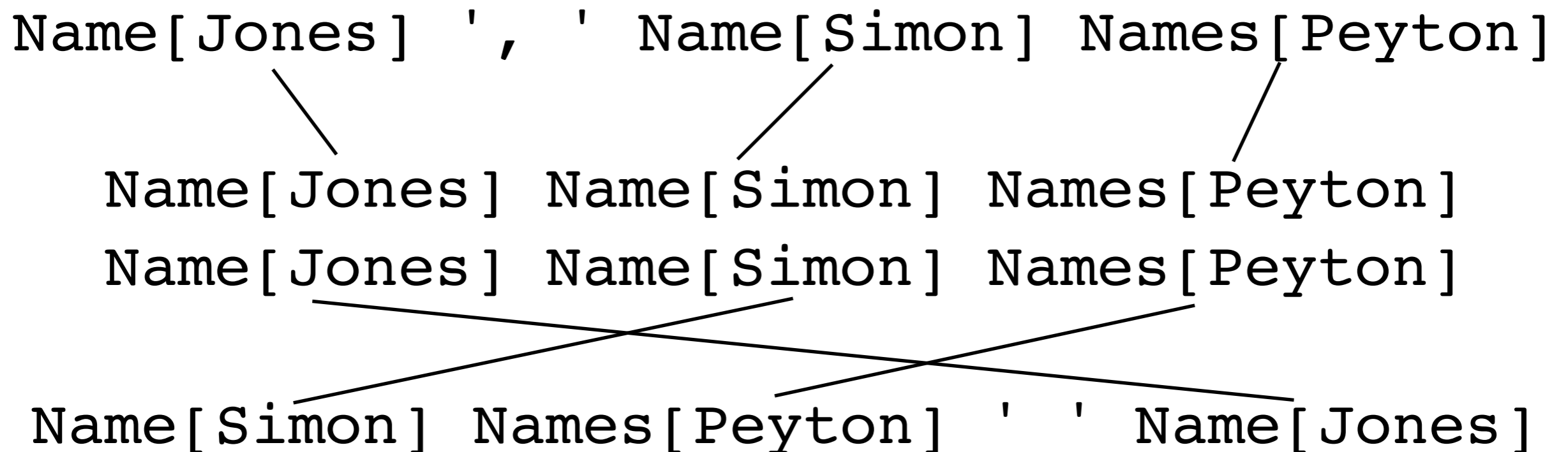
```
Name [ Jones ] Name [ Simon ] Names [ Peyton ]
```

```
Name [ Jones ] Name [ Simon ] Names [ Peyton ]
```

```
Name [ Simon ] Names [ Peyton ] ' ' Name [ Jones ]
```

Satisfying Examples

Add in additional information from the parsing of input/
output example strings!



Satisfying Examples

Add in additional information from the parsing of input/
output example strings!

~~Name [Jones] ' , ' Name [Simon] Names [Peyton]~~
Name [Jones] Name [Simon] Names [Peyton]
Name [Jones] Name [Simon] Names [Peyton]
~~Name [Simon] Names [Peyton] ' ' Name [Jones]~~

Evaluation

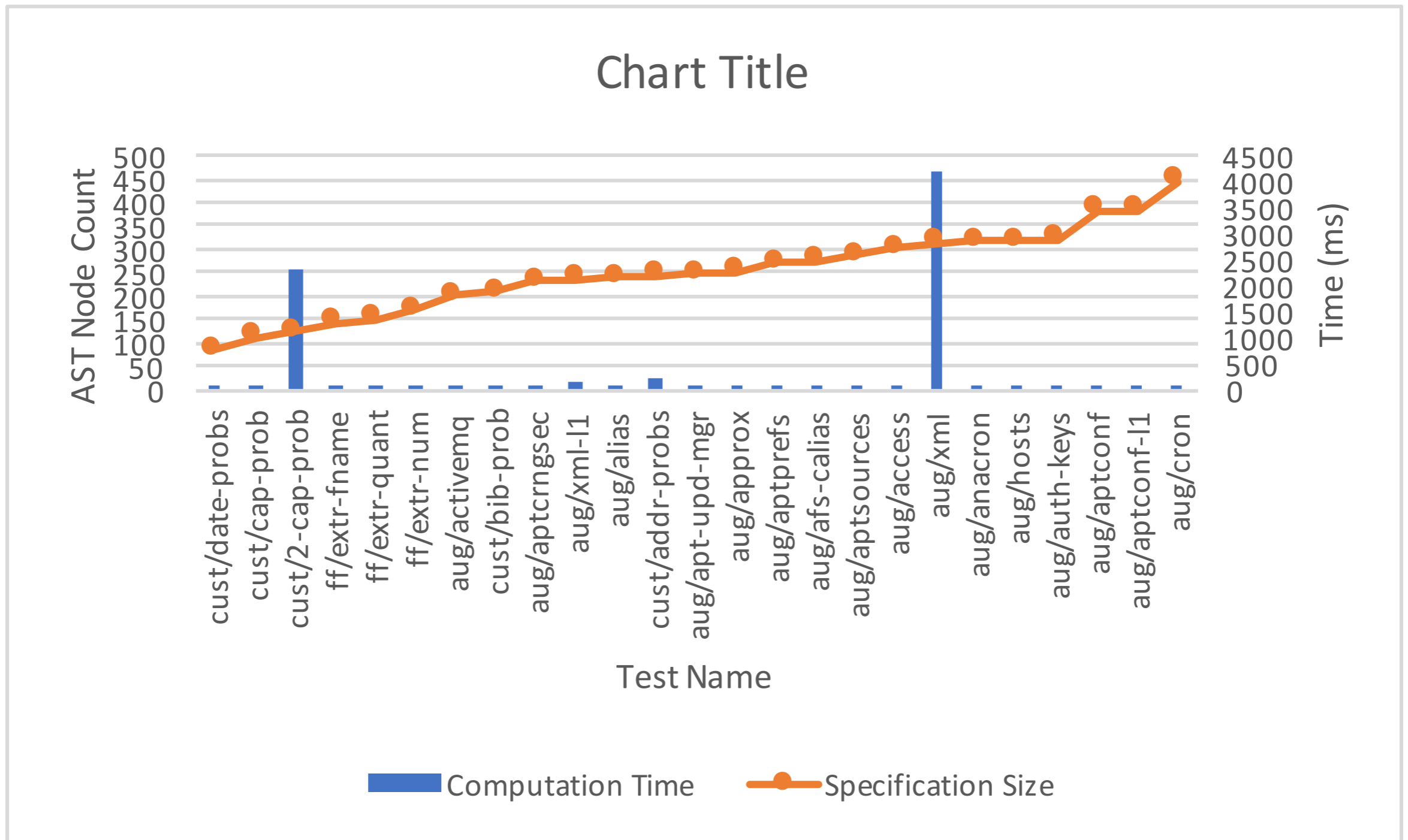
Evaluation Tests

- Custom built
 - Examples of good uses for it
 - Adversarial Examples
- FlashFill
 - Related work on string transformations
- Augeas
 - Bidirectional Language between configuration files and tree representation of the information

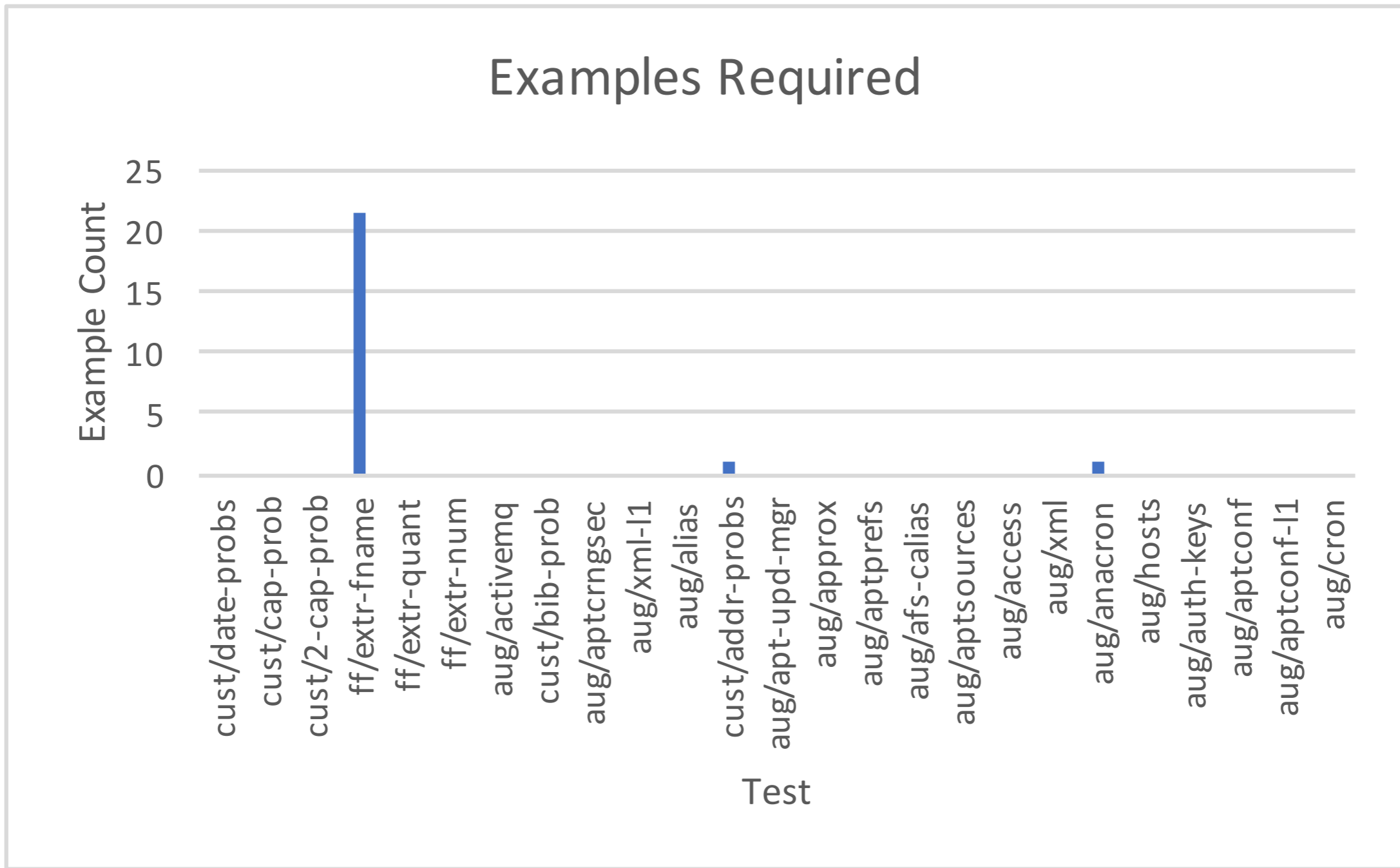
Evaluation Methods

- Evaluated on this very Mac!
 - 2.5 GHz Intel Core i7 processor
 - 16 GB of 1600 MHz DDR3
 - OS X Yosemite
- Average of 10 runs

Synthesis Time Graph



Examples Graph



What was hard?

- MANY were easy (13/25 synthesized in <10ms)
- custom 2-cap-prob (2.3s)
 - Exploits the fact we always distribute
 - Analogous 4-cap-prob takes 30 minutes
- Augeas's xml
 - Source and Target are very different
 - Large number of conversions needed to take place
 - 12 conversions required
 - 400 expansions explored

Surprising Facts

- Examples were often unnecessary
 - Tight type system makes few lenses per type
 - Usually little reordering took place
- How well it scales in practice!
 - Usually more complex programs become quickly impossible to synthesize
 - Difficulty comes from how different the formats are, not how complex the formats were

Related Work

- Type-Directed Synthesis [OZ '15] [FOWZ '16] [FCD '15] [GKKP '13] ...
 - In general purpose programming languages
 - Type systems very different
- String Transformation Synthesis [G '11] [LG '14] [PGDP '14] [S '16] [SG '12] ...
 - Did not take in specifications and examples

Synthesizing Bijective String Transformers

- Developed **DNF Lenses**
- Proven DNF Lenses **sound and complete** with respect to Bijective Lenses
- Developed an **efficient algorithm** for synthesis of DNF lenses
- **Evaluated our algorithm** on a suite of examples consisting of our own examples, examples derived from related work, and real world examples